

ZERO FORCING NUMBER, MAXIMUM NULLITY, AND PATH COVER NUMBER OF SUBDIVIDED GRAPHS*

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1 **Abstract.** The zero forcing number, maximum nullity and path cover number of a (simple, undirected) graph
2 are parameters that are important in the study of minimum rank problems. We investigate the effects on these
3 graph parameters when an edge is subdivided to obtain a so-called edge subdivision graph. An open question raised
4 by Barrett et al. in “Minimum rank of edge subdivisions of graphs,” *Electronic Journal of Linear Algebra* (2009) 18:
5 530–563, is answered in the negative, and we provide additional evidence for an affirmative answer to another open
6 question in that paper. It is shown that there is an independent relationship between the change in maximum nullity
7 and zero forcing number caused by subdividing an edge once. Bounds on the effect of a single edge subdivision on
8 the path cover number are presented, conditions under which the path cover number is preserved are given, and it is
9 shown that the path cover number and the zero forcing number of a complete subdivision graph need not be equal.

10 **Keywords.** zero forcing number, maximum nullity, minimum rank, path cover number, edge
11 subdivision, matrix, multigraph, graph

12 **AMS subject classifications.** 05C50, 15A03, 15A18, 15B57

13 **1. Introduction.** Let F be any field. For a (simple, undirected) graph $G = (V, E)$ that has
14 vertex set $V = \{1, \dots, n\}$ and edge set E , $\mathcal{S}(F, G)$ is the set of all symmetric $n \times n$ matrices A
15 with entries from F such that for any non-diagonal entry a_{ij} in A , $a_{ij} \neq 0$ if and only if $ij \in E$.
16 The *minimum rank* of G is

$$17 \quad \text{mr}(F, G) = \min\{\text{rank } A : A \in \mathcal{S}(F, G)\},$$

18 and the *maximum nullity* of G is

$$19 \quad \text{M}(F, G) = \max\{\text{null } A : A \in \mathcal{S}(F, G)\}.$$

20 Note that $\text{mr}(F, G) + \text{M}(F, G) = |G|$, where $|G|$ is the number of vertices in G . Thus the problem
21 of finding the minimum rank of a given graph is equivalent to the problem of determining its
22 maximum nullity.

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23 We say that a graph $H = (V', E')$ is a *subgraph* of $G = (V, E)$ if $V' \subseteq V$ and $E' \subseteq E$. The
 24 subgraph H is called an *induced subgraph* if for each $x, y \in V', xy \in E'$ if and only if $xy \in E$.
 25 Denote by $G[X]$ the induced subgraph of G with vertex set $X \subseteq V$; $G - W$ is used to denote
 26 $G[V \setminus W]$. The graph $G - \{v\}$ is also denoted by $G - v$. A graph G is the *union* of graphs
 27 $G_i = (V_i, E_i)$, $1 \leq i \leq h$, if $G = (\cup_{i=1}^h V_i, \cup_{i=1}^h E_i)$. A vertex v of a connected graph G is a
 28 *cut-vertex* if $G - v$ is disconnected. An edge e of a connected graph G is a *cut-edge* if $G - e$ is
 29 disconnected. The *rank spread* of G is $r_v(F, G) = \text{mr}(F, G) - \text{mr}(F, G - v)$. One technique in
 30 computing minimum rank is by *cut-vertex reduction* (see, e.g., [6]), which is as follows: Suppose
 31 that v is a cut-vertex of G . For $i = 1, \dots, h$, let $W_i \subseteq V(G)$ be the vertices of the i th component of
 32 $G - v$ and let $G_i = G[\{v\} \cup W_i]$. Then $\text{mr}(F, G) = \sum_{i=1}^h \text{mr}(F, G_i - v) + \min\{2, \sum_{i=1}^h r_v(F, G_i)\}$.
 33 For a graph $G = (V, E)$, the *degree* of $v \in V$, denoted $\text{deg } v$, is the number of vertices in V that
 34 share an edge with v . A *leaf* vertex is a vertex of degree one. A *high degree* vertex is a vertex of
 35 degree greater than or equal to 3.

36 OBSERVATION 1.1. *Let G be a graph, let v be a leaf vertex of a graph G , and let F be a field.*
 37 *It is easy to see that $\text{mr}(F, G) - \text{mr}(F, G - v) \leq 1$, or equivalently, $M(F, G) \geq M(F, G - v)$.*

38 We consider two graph parameters that are related to the maximum nullity, namely the zero
 39 forcing number and the path cover number. The zero forcing number of a graph is the minimum
 40 number of black vertices initially needed to color all vertices black according to the color-change
 41 rule. The *color-change rule* is defined as follows: if G is a graph with each vertex colored either
 42 white or black, u is a black vertex of G and exactly one neighbor v of u is white, then change the
 43 color of v to be black. Let S be a subset of V . The *derived coloring of S* is the result of coloring
 44 every vertex in S black and every vertex not in S white, and then applying the color-change rule
 45 until no more changes are possible. A *zero forcing set* of G is a set $Z \subseteq V$ such that every vertex
 46 in the derived coloring of Z is black. The *zero forcing number* of G is

$$47 \quad Z(G) = \min\{|Z| : Z \text{ is a zero forcing set of } G\}.$$

48 A zero forcing set of G , Z , is called a *minimum zero forcing set* of G if $|Z| = Z(G)$.

49 A *path* in G is a subgraph $H = (V', E')$ where $V' = \{u_1, \dots, u_k\}$ and
 50 $E' = \{u_1u_2, u_2u_3, \dots, u_{k-1}u_k\}$; a path is *even* or *odd* according as its number of vertices is even or
 51 odd. A *Hamiltonian path* of a graph G is a path that includes all the vertices of G . A *path cover*
 52 of G is a set of vertex disjoint paths, each of which is an induced subgraph of G , that contains all
 53 vertices of G . The *path cover number* of G is

$$54 \quad P(G) = \min\{|\mathcal{P}| : \mathcal{P} \text{ is a path cover of } G\}.$$

55 A path cover of G , \mathcal{P} , is called a *minimum path cover* of G if $|\mathcal{P}| = P(G)$.

56 The relationships between $M(F, G)$, $Z(G)$ and $P(G)$ for any graph G are discussed in papers
 57 devoted to the study of minimum rank problems. For extensive surveys on minimum rank and
 58 related problems, see [6] or [7].

59 THEOREM 1.2. [1] *For any graph G , $M(F, G) \leq Z(G)$.*

60 THEOREM 1.3. [8] *For any graph G , $P(G) \leq Z(G)$.*

61 In [2], examples of graphs are given to show that both $M(F, G) < P(G)$ and $P(G) < M(F, G)$

62 are possible. In particular, $M(F, G) < Z(G)$ is possible. However, all three parameters give equality
 63 for graphs that are trees.

64 THEOREM 1.4. [1, 5, 9] *For any tree T , $M(F, T) = P(T) = Z(T)$.*

65 Following the notation in [3], we give the following definitions. Let $e = uv$ be an edge of G .
 66 Define G_e to be the graph obtained from G by inserting a new vertex w into V , deleting the edge
 67 e and inserting edges uw and wv . We say that the edge e has been *subdivided* and call G_e an
 68 *edge subdivision* of G . A *complete subdivision graph* \vec{G} is obtained from a graph G by subdividing
 69 every edge of G once. In [3] and [10], the authors investigate the maximum nullity and zero forcing
 70 number of graphs obtained by a finite number of edge subdivisions of a given graph and, among
 71 other results, establish the following two propositions about the effect of an edge subdivision on
 72 the zero forcing number and maximum nullity.

73 PROPOSITION 1.5. [3, 10] *Let G be a graph and let e be an edge of G . Then*

74
$$M(F, G) \leq M(F, G_e) \leq M(F, G) + 1 \quad \text{and} \quad Z(G) \leq Z(G_e) \leq Z(G) + 1.$$

75 PROPOSITION 1.6. [3, 10] *Let G be a graph and let e be an edge of G incident to a vertex of
 76 degree at most 2. If $F \neq \mathbb{Z}_2$, then $M(F, G) = M(F, G_e)$ and $Z(G) = Z(G_e)$.*

78 The paper [3] concludes with a list of open questions, including the following two questions.

79 QUESTION 1.7. *Let F be a field. Suppose G is a graph in which each vertex has degree at
 80 least 3 and H is a graph that has one less edge subdivision than \vec{G} . Is it always the case that
 81 $M(F, H) < M(F, \vec{G})$?*

82 QUESTION 1.8. *Is $M(F, \vec{G}) = Z(\vec{G})$ for every field F and graph G ?*

83 In [3], the authors provide the following substantial result toward an affirmative answer to
 84 Question 1.8.

85 THEOREM 1.9. [3] *If $G = (V, E)$ has a Hamiltonian path then $M(F, \vec{G}) = Z(\vec{G}) = m - n + 2$
 86 and $\text{mr}(F, \vec{G}) = 2n - 2$, where $n = |V|$ and $m = |E|$.*

87 In Section 2 we provide additional evidence of an affirmative answer to Question 1.8, including
 88 establishing that $M(F, \vec{G}) = Z(\vec{G})$ if G does not have a cut-edge. In Section 3 we give an example
 89 that provides a negative answer to Question 1.7. We also present examples showing that there
 90 is an independent relationship between the change in maximum nullity and zero forcing number
 91 caused by a single edge subdivision in a graph G . In Section 4, we give bounds on the effect of a
 92 single edge subdivision on the path cover number and give conditions under which the path cover
 93 number is preserved. We also provide an example to show that $P(\vec{G})$ need not equal $Z(\vec{G})$ for an
 94 arbitrary graph G .

95 **2. Complete edge subdivision graphs.** In [3] it was shown that $M(F, \vec{G}) = Z(\vec{G})$ if G
 96 has a Hamiltonian path. In this section we establish $M(F, \vec{G}) = Z(\vec{G})$ for other conditions on G ,
 97 specifically for graphs G such that G is a cactus or has no cut-edge.

98 A *cactus* is a graph in which any two cycles share at most one vertex. We use Row's work on
 99 cacti to show that the zero forcing number and maximum nullity of a complete subdivision of any

100 cactus is equal.

101 PROPOSITION 2.1. [11] *Let G be a cactus in which each cycle has three vertices, an even*
102 *number of vertices, or a vertex which has only two neighbors. Then $M(\mathbb{R}, G) = Z(G)$.*

103 PROPOSITION 2.2. *If $G = (V, E)$ is a cactus, then $M(F, \vec{G}) = Z(\vec{G})$.*

104 *Proof.* Let $G = (V, E)$ be a cactus. We perform a complete subdivision on G . Notice then
105 that \vec{G} is a cactus. Furthermore, each cycle in \vec{G} is even (and has a vertex of degree two). Thus
106 $M(\mathbb{R}, \vec{G}) = Z(\vec{G})$. If H is a cycle or tree, then $M(F, H) = M(\mathbb{R}, H)$. Since cut-vertex reduction
107 preserves field independence (see [6]), $M(F, \vec{G}) = Z(\vec{G})$ for every cactus G . \square

108 To prove that $M(F, \vec{G}) = Z(\vec{G})$ for every G that does not have a cut-edge, we first generalize
109 the set of complete edge subdivision graphs.

110 DEFINITION 2.3. Let \mathcal{K} be the family of bipartite graphs $G = (V(G), E(G))$ such that there
111 is a bipartition of the vertices $V(G) = X \dot{\cup} Y$ with $\deg x \leq 2$ for all $x \in X$.

112 Note that every path is in \mathcal{K} , and every even cycle is in \mathcal{K} . An odd cycle is not bipartite, so
113 it is not in \mathcal{K} . If G is any connected bipartite graph, then the (unordered) pair of bipartition sets
114 is uniquely determined. If $G \in \mathcal{K}$ and G has a high degree vertex, then the bipartition sets X and
115 Y such that $V(G) = X \dot{\cup} Y$ and $\deg x \leq 2$ for all $x \in X$ are uniquely determined. When the sets
116 X, Y such that $V(G) = X \dot{\cup} Y$ and $\deg x \leq 2$ for all $x \in X$ are not uniquely determined, we often
117 make a choice, possibly subject to some additional condition(s). When X and Y are specified by
118 uniqueness or by choice, we write $X(G)$ for X and $Y(G)$ for Y .

119 PROPOSITION 2.4. *A graph H is a complete subdivision graph of some graph G if and only if*
120 *$H \in \mathcal{K}$, H does not contain a cycle on four vertices, and $\deg x = 2$ for every $x \in X(H)$.*

121 *Proof.* The forward direction is clear. For the converse, we reconstruct G from H . It is
122 sufficient to do so for a connected graph, and then take the union of connected components, so
123 assume H is connected. If H has no high degree vertex, then H is an even cycle or odd path (an
124 even path is not allowed because one vertex in each bipartition set of such a path has degree one),
125 and thus H is a complete subdivision graph. So assume H has a high degree vertex. For each
126 $x \in X(H)$ with neighbors $y_1, y_2 \in Y(H)$, delete edges xy_1 and xy_2 and vertex x and add edge
127 y_1y_2 . This method creates a graph G such that $H = \vec{G}$: G is a graph, since no duplicate edges
128 are created (two vertices $x_1, x_2 \in X$ with the same neighbors $y_1, y_2 \in Y(G)$ would have created a
129 cycle on four vertices in H , which we expressly disallow). \square

130 CONJECTURE 2.5. *If $G \in \mathcal{K}$, then $M(F, G) = Z(G)$.*

131 By Proposition 2.4, every complete subdivision graph is in \mathcal{K} , so this conjecture generalizes a
132 conjecture that $M(F, \vec{G}) = Z(\vec{G})$ for all graphs G .

133 The method by which we show $M(F, \vec{G}) = Z(\vec{G})$ for graphs without a cut-edge requires knowing
134 that certain diagonal entries of a matrix are zero. A graph $G \in \mathcal{K}$ is *special* if there exists a matrix
135 $A \in \mathcal{S}(G)$ such that

- 136 1. null $A = M(F, G)$.
- 137 2. If $x \in X(G)$, then $a_{xx} = 0$.

138 For a special graph G , a matrix $A \in \mathcal{S}(G)$ satisfying conditions (1) and (2) is *optimal* for G .

139 Let G be a graph and let $C = (V_C, E_C)$ be a cycle that is a subgraph of G . A *subdivided*
 140 *chordal path* of G is a path $P = (v_1, \dots, v_{2k+1})$ in G such that $v_1, v_{2k+1} \in V_C$, $\deg_G v_i = 2$ for
 141 $i = 2, 3, \dots, 2k$, and $v_i \notin V_C$ for $i = 2, 3, \dots, 2k$.

142 **THEOREM 2.6.** *Let G' be a graph in \mathcal{K} and let G be obtained from G' by removing a subdivided*
 143 *chordal path $P = (v_1, v_2, v_3)$ of G' between two vertices in $V(G)$. If $M(F, G) = Z(G)$ and G is*
 144 *special, then $M(F, G') = Z(G')$ and G' is special.*

145 *Proof.* Suppose that $M(F, G) = Z(G)$ and G is special. Let $Q = (v_1, u_2, \dots, u_{2k}, v_3)$ be another
 146 path that connects v_1 and v_3 . Since $G' \in \mathcal{K}$ and $v_1, v_3 \in Y(G')$, $\deg_G u_{2i} = \deg_{G'} u_{2i} = 2$ for
 147 $i = 1, \dots, k$. Let A be an optimal matrix for G , so the diagonal entries of A in the column vectors
 148 $\mathbf{a}_{u_{2i}}$ associated with vertices $u_{2i}, i = 1, \dots, k$ are all zero. Since the only vertices adjacent to u_2 are
 149 v_1 and u_3 , \mathbf{a}_{u_2} has nonzero entries exactly in rows v_1 and u_3 , and similarly, \mathbf{a}_{u_4} has nonzero entries
 150 exactly in rows u_3 and u_5 . We can take a linear combination of these two vectors to cancel the
 151 nonzero entry in row u_3 , to obtain a column vector with nonzero entries exactly in rows v_1, u_5 . We
 152 iterate this process with column vectors to obtain a column vector \mathbf{c} with non-zero entries in exactly
 153 rows v_1, v_3 . Let $A' = [a'_{ij}]$ be A with the extra column \mathbf{c} and extra row \mathbf{c}^T and zero as the new
 154 diagonal entry. We know $A' \in \mathcal{S}(G')$. Since G is an induced subgraph of G' , $\text{mr}(F, G) \leq \text{mr}(F, G')$.
 155 Since $\text{rank}(A') = \text{rank}(A)$, $\text{mr}(F, G) = \text{mr}(F, G')$. Hence, $M(F, G') = M(F, G) + 1$.

156 Since $a'_{xx} = 0$ for every $x \in X(G')$, G' is special. Note that $Z(G) + 1 = M(F, G) + 1 =$
 157 $M(F, G') \leq Z(G') \leq Z(G) + 1$. Hence, $Z(G') = M(F, G')$. \square

158 Although this paper is primarily concerned with simple graphs, multigraphs are a useful tool.
 159 A *multigraph* $G = (V, E)$ is a general graph in which E is a multiset of two-element subsets of
 160 vertices. That is, a multigraph allows multiple copies of an edge vw (where $v \neq w$), but a loop vv
 161 is not permitted. For a field $F \neq \mathbb{Z}_2$, the *maximum nullity of a multigraph* G of order n over F ,
 162 denoted by $M(F, G)$, is the largest possible nullity over all matrices $A \in F^{n \times n}$ whose ij th entry
 163 a_{ij} (for $i \neq j$) is zero if i and j are not adjacent in G , is nonzero if ij is a single edge, and is any
 164 element of F if ij is a multiple edge. In the case that $F = \mathbb{Z}_2$ and ij is a multiple edge, a_{ij} is
 165 0 if the number of copies of edge ij is even and 1 if it is odd. If a multigraph does not have any
 166 multiple edges then it is a (simple) graph. Observe that if G is a multigraph, then \overline{G} is a (simple)
 167 graph and $\overline{G} \in \mathcal{K}$.

168 The *contraction* of edge $e = uv$ of G is the multigraph obtained from G by identifying the
 169 vertices u and v , deleting any loops that arise in this process. A set $R \subset V(G)$ is a *separating set*
 170 of a graph G if $G - R$ has more connected components than G does; in this case R is called an
 171 r -separating set where $r = |R|$. A 1-separating set is a cut-vertex, and cut-vertex reduction is a
 172 standard technique for computing minimum rank/maximum nullity. Van der Holst has established
 173 a 2-separating set reduction for computing maximum nullity using multigraphs. A *2-separation*
 174 of G is a pair of subgraphs $(G_1(R), G_2(R))$ such that $V(G_1(R)) \cap V(G_2(R)) = R = \{r_1, r_2\}$,
 175 $V(G_1(R)) \cup V(G_2(R)) = V(G)$, $E(G_1(R)) \cap E(G_2(R)) = \emptyset$, and $E(G_1(R)) \cup E(G_2(R)) = E(G)$.
 176 We introduce some notation for the multigraphs needed for van der Holst's 2-separation theorem.
 177 For $i = 1, 2$, $H_i(R)$ is the graph or multigraph obtained from $G_i(R)$ by adding edge $r_1 r_2$. If
 178 $r_1 r_2 \notin E(G_i(R))$, $H_i(R)$ is a (simple) graph; otherwise $H_i(R)$ is a multigraph having two edges

179 between r_1 and r_2 (with every other pair of vertices either nonadjacent or joined by exactly one
 180 edge). At most one of $H_1(R), H_2(R)$ has a multiple edge. For $i = 1, 2$, $\widehat{G}_i(R)$ is the multigraph
 181 obtained from $H_i(R)$ by contracting an edge r_1r_2 (note that van der Holst uses the notation $\overline{G}_i(R)$
 182 for what we denote by $\widehat{G}_i(R)$, but $\overline{G}_i(R)$ may cause confusion with a complement).

183 THEOREM 2.7. [12] Let G be a (simple) graph, let $(G_1(R), G_2(R))$ be a 2-separation of G .
 184 Then

$$185 \quad M(F, G) = \max \left\{ \begin{array}{l} M(F, G_1(R)) + M(F, G_2(R)), \\ M(F, H_1(R)) + M(F, H_2(R)), \\ M(F, \widehat{G}_1(R)) + M(F, \widehat{G}_2(R)), \\ M(F, G_1(R) - r_1) + M(F, G_2(R) - r_1), \\ M(F, G_1(R) - r_2) + M(F, G_2(R) - r_2), \\ M(F, G_1(R) - R) + M(F, G_2(R) - R) \end{array} \right\} - 2.$$

186
 187 LEMMA 2.8. Let G be a graph in \mathcal{K} and $(G_1(R), G_2(R))$ be a 2-separation of G . If $G_1(R)$ is an
 188 even path with endpoints r_1 and r_2 and $r_1r_2 \notin E(G)$, then $M(F, G) = M(F, H_1(R)) + M(F, H_2(R)) - 2$
 189 (or equivalently, $\text{mr}(F, G) = \text{mr}(F, H_1(R)) + \text{mr}(F, H_2(R))$) and $H_1(R), H_2(R) \in \mathcal{K}$.

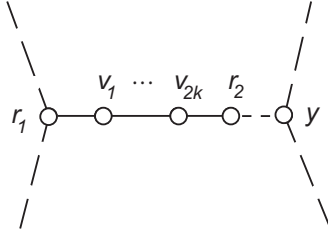


Fig. 2.1: Illustration for Lemma 2.8

190 *Proof.* Let $G_i = G_i(R), H_i = H_i(R), \widehat{G}_i = \widehat{G}_i(R), i = 1, 2$. Since $r_1r_2 \notin E(G)$, H_1 and H_2 are
 191 (simple) graphs, and it is clear that $H_1, H_2 \in \mathcal{K}$. To show $M(F, G) = M(F, H_1) + M(F, H_2) - 2$,
 192 by Theorem 2.7 it suffices to prove the following inequalities.

- 193 • $M(F, H_1) + M(F, H_2) \geq M(F, G_1) + M(F, G_2)$: Since G_1 is a path and H_1 is a cycle,
 194 $M(F, G_1) = M(F, H_1) - 1$. Since G_2 is obtained from H_2 by deleting the edge r_1r_2 ,
 195 $M(F, H_2) \geq M(F, G_2) - 1$. Hence,

$$196 \quad M(F, H_1) + M(F, H_2) \geq M(F, G_1) + 1 + M(F, G_2) - 1 \\
 197 \quad = M(F, G_1) + M(F, G_2).$$

- 198 • $M(F, H_1) + M(F, H_2) \geq M(F, \widehat{G}_1) + M(F, \widehat{G}_2)$: Since \widehat{G}_1 is a cycle, $M(F, \widehat{G}_1) = 2 =$
 199 $M(F, H_1)$. If $\deg r_2 = 1$, then r_2 is a leaf of H_2 , so by Observation 1.1, $M(F, H_2) \geq$
 200 $M(F, H_2 - r_2) = M(F, \widehat{G}_2)$. So assume $\deg r_2 = 2$ and let $r_2y \in E(G)$ and $y \neq v_{2k}$. Note
 201 that $r_1y \notin E(G)$ since r_1, y are in the same bipartition set and $r_1 \neq y$. Observe that
 202 $H_2 = (\widehat{G}_2)_e$ where $e = r_2y$. By Proposition 1.5, $M(F, \widehat{G}_2) \leq M(F, H_2)$, and the desired
 203 inequality follows.

- 204 • For $i = 1, 2$, $M(F, H_1) + M(F, H_2) \geq M(F, G_1 - r_i) + M(F, G_2 - r_i)$: Observe that $M(F, G_1 - r_i) = 1 = M(F, H_1) - 1$. Since $G_2 - r_i = H_2 - r_i$, $M(F, H_2) \geq M(F, H_2 - r_i) - 1 = M(F, G_2 - r_i) - 1$, and the desired inequality follows.
- 205
- 206 • $M(F, H_1) + M(F, H_2) \geq M(F, G_1 - R) + M(F, G_2 - R)$: Observe that $M(F, G_1 - R) = 1 = M(F, H_1) - 1$. Since $G_2 - r_1 = H_2 - r_1$, $M(F, H_2) \geq M(F, H_2 - r_1) - 1 = M(F, G_2 - r_1) - 1$. Since r_2 is a leaf vertex of $G_2 - r_1$, $M(F, G_2 - R) \leq M(F, G_2 - r_1)$, and thus $M(F, H_2) \geq M(F, G_2 - R) - 1$. Hence the desired inequality follows.
- 207
- 208
- 209
- 210

□

211 If $V(L) \subset V(G)$ and $A = [a_{uv}] \in \mathcal{S}(L)$, then the *embedding* $\tilde{A} = [\tilde{a}_{uv}]$ of A for G is the $|G| \times |G|$
 212 matrix defined by $\tilde{a}_{uv} = a_{uv}$ if $u, v \in V(L)$ and 0 otherwise. A *decomposition* of a graph G is a
 213 pair of graphs (L_1, L_2) such that

- 214 1. $V(G) = V(L_1) \cup V(L_2)$.
- 215 2. $|V(L_1) \cap V(L_2)| = 2$.
- 216 3. $|E(L_1) \cap E(L_2)| = 0$ or 1.
- 217 4. $E(G) = (E(G_1) \cup E(G_2)) \setminus (E(G_1) \cap E(G_2))$.

218 Every 2-separation $(G_1(R), G_2(R))$ of G is a decomposition of G , but not conversely. A decom-
 219 position (L_1, L_2) of a graph $G \in \mathcal{K}$ is a *special decomposition* if it satisfies all of the following
 220 conditions:

- 221 1. $L_1, L_2 \in \mathcal{K}$.
- 222 2. $\text{mr}(F, G) = \text{mr}(F, L_1) + \text{mr}(F, L_2)$. Equivalently, $M(F, G) = M(F, L_1) + M(F, L_2) - 2$.
- 223 3. For $r \in V(L_1) \cap V(L_2)$, either $r \in Y(L_1) \cap Y(L_2)$ or $r \in X(L_1) \cap X(L_2)$.

224 LEMMA 2.9. *Suppose (L_1, L_2) is a decomposition of a graph G . If $A_k \in \mathcal{S}(L_k), k = 1, 2$, then*
 225 *there exists $\alpha \in F$ such that $A = A_1 + \alpha A_2 \in \mathcal{S}(G)$. If $\text{mr}(F, G) = \text{mr}(F, L_1) + \text{mr}(F, L_2)$ and*
 226 *rank $A_k = \text{mr}(F, L_k)$, for $k = 1, 2$, then rank $A = \text{mr}(F, G)$ (for this α). If (L_1, L_2) is a special*
 227 *decomposition of $G \in \mathcal{K}$ and L_1 and L_2 are special, then G is special.*

228 *Proof.* If $E(L_1) \cap E(L_2) = \emptyset$, choose $\alpha = 1$. If $E(L_1) \cap E(L_2) = \{zw\}$ choose $\alpha = -a_{zw}^{(1)}/a_{zw}^{(2)}$
 229 where $A_k = [a_{ij}^{(k)}], k = 1, 2$. Then $A \in \mathcal{S}(G)$ and $\text{rank } A \leq \text{rank } A_1 + \text{rank } A_2$, so $\text{mr}(F, G) =$
 230 $\text{mr}(F, L_1) + \text{mr}(F, L_2)$ implies $\text{rank } A = \text{mr}(F, G)$.

231 Now suppose (L_1, L_2) is a special decomposition of G and L_1, L_2 are special. Construct
 232 $A = [a_{ij}]$ as previously using optimal A_k for $L_k, k = 1, 2$. We claim A is optimal for G and thus G
 233 is special. It is already established that $\text{null } A = M(F, G)$ and since for $r \in V(L_1) \cap V(L_2)$, either
 234 $r \in Y(L_1) \cap Y(L_2)$ or $r \in X(L_1) \cap X(L_2)$, the required zeros on the diagonal are preserved. □

235 THEOREM 2.10. *Let G' be a graph in \mathcal{K} and let G be obtained from G' by removing a subdivided*
 236 *chordal path $P = (v_1, \dots, v_{2k+1})$ of G' between two vertices in $V(G)$. If $M(F, G) = Z(G)$ and G is*
 237 *special, then $M(F, G') = Z(G')$ and G' is special.*

238 *Proof.* Theorem 2.6 covers the case $k = 1$, so assume $k \geq 2$. Let $r_1 = v_1, r_2 = v_{2k}$, and
 239 $R = \{r_1, r_2\}$. Let $G_1(R) = (r_1, v_2, \dots, v_{2k-1}, r_2)$ be a path in G' and $G_2(R) = G' - \{v_2, \dots, v_{2k-1}\}$,
 240 so $(G_1(R), G_2(R))$ is a 2-separation of G' ; see Figure 2.2. Since $r_1 r_2 \notin E(G')$, H_1 is a cycle on
 241 $2k$ vertices and H_2 is obtained from G by adding the subdivided chordal path (v_1, r_2, v_{2k+1}) ; see

242 Figure 2.2. By Theorem 2.6 H_2 is special, and by Lemma 2.8 $\text{mr}(F, G') = \text{mr}(F, H_1) + \text{mr}(F, H_2)$.
 243 Thus (H_1, H_2) is a special decomposition of G' , and so by Lemma 2.9, G' is special. Furthermore,
 244 we have

$$\begin{aligned}
 245 \quad M(F, G') &= M(F, H_1) + M(F, H_2) - 2 \\
 246 \quad &= M(F, H_2) \\
 247 \quad &= Z(H_2) \\
 248 \quad &= Z(G')
 \end{aligned}$$

249 by subdividing edges incident to a vertex of degree two. \square

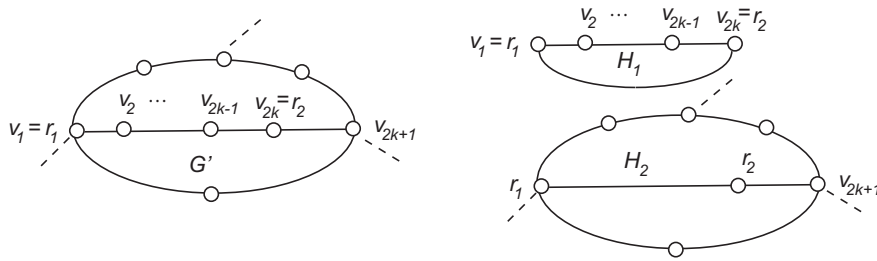


Fig. 2.2: Illustration for Theorem 2.10

250 LEMMA 2.11. Let G be a graph. If cycles C_1, C_2 of G intersect in $k > 1$ paths, then there is
 251 a cycle C_3 of G such that C_1 and C_3 intersect in exactly one path and that path has at least two
 vertices.

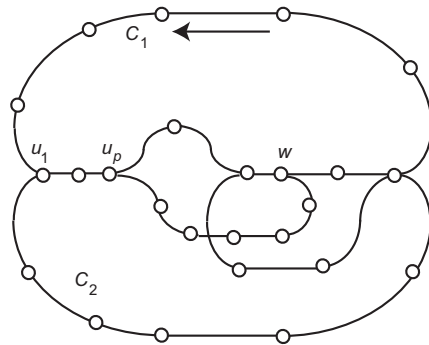


Fig. 2.3: Illustration for Lemma 2.11

252 *Proof.* Choose an orientation for C_1 . With this orientation, each vertex $v \in C_1$ has a prede-
 253 cessor and a successor. Let $P = (u_1, \dots, u_p)$ be a path in $C_1 \cap C_2$ that conforms to the orientation
 254 and that is maximal in the sense that the predecessor of u_1 in C_1 is not in C_2 and the successor
 255 of u_p in C_1 is not in C_2 . Impose the orientation of P on C_2 . Let w be the first vertex in C_2 after
 256 u_p that is also in C_1 (see Figure 2.3). Let P_i be the path in C_i connecting u_p and w (following
 257 the orientation of C_i). Define C_3 to be the cycle enclosed by P_1 and P_2 . Then C_1 intersects C_3 in
 258 exactly P_1 , and $u_p, w \in V(P_1)$. \square

260 LEMMA 2.12. Let G be a graph in \mathcal{K} . Suppose cycles C_1, C_2 of G intersect in exactly one path
 261 P and none of the interior vertices of P is a cut-vertex. Then G contains a subdivided chordal
 262 path of some cycle.

263 *Proof.* Let $P = (v_1, \dots, v_m)$. The proof is by strong induction on the number ℓ of high degree
 264 vertices among the interior vertices $v_i, i = 2, \dots, m - 1$. If $\ell = 0$, then P is a subdivided chordal
 265 path of G . So assume that if two cycles of G intersect in exactly one path that has $q < \ell$ high
 266 degree interior vertices, then G contains a subdivided chordal path, and suppose P has ℓ high
 267 degree interior vertices. Let v_t be a high degree interior vertex. Since v_t is not a cut-vertex, there
 268 exists a path Q_1 that connects v_t to some other vertex $y \in V(C_1)$ (if necessary reverse the names
 269 of C_1 and C_2) and such that $V(Q) \cap V(C_1) = \{v_t, y\}$. We consider two cases depending on whether
 270 or not y is on P , as illustrated in Figure 2.4.

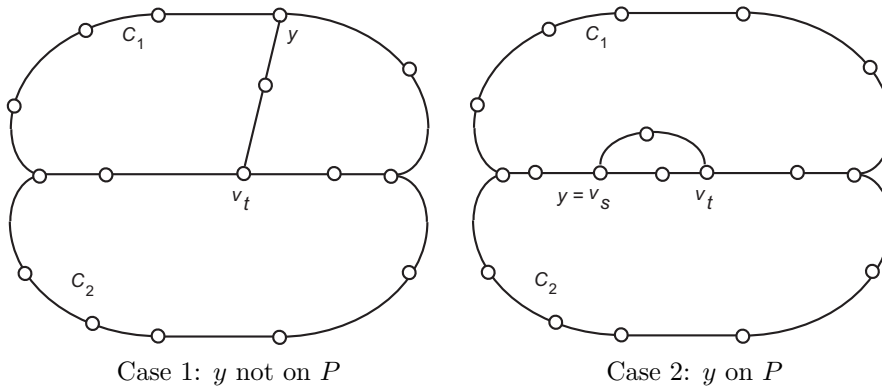


Fig. 2.4: Illustration for Lemma 2.12

271 **Case 1.** $y \notin V(P)$: Let Q_2 be the path in C_1 between y and v_t that does not contain v_m .
 272 Then $(v_1, v_2, \dots, v_t), Q_1$, and Q_2 form a cycle C_3 that intersects C_2 in path $P' = (v_1, v_2, \dots, v_t)$.
 273 Since P' has fewer high degree interior vertices, G contains a subdivided chordal path.

274 **Case 2.** $y \in V(P)$: Let P' be the subpath of P between $v_s = y$ and v_t , so P' and Q_1 form
 275 a cycle C_3 that intersects C_2 in path $P' = (v_s, \dots, v_t)$. Since P' has fewer high degree interior
 276 vertices, G contains a subdivided chordal path. \square

277 PROPOSITION 2.13. Suppose G has a cut-vertex v . For $i = 1, \dots, h$, let $W_i \subseteq V(G)$ be
 278 the vertices of the i th component of $G - v$ and let G_i be the subgraph induced by $\{v\} \cup W_i$. If
 279 $r_v(F, G_1) = 0$, then

$$\text{mr}(F, G) = \text{mr}(F, G_1) + \text{mr}(F, G - W_1).$$

280
 281
 282 *Proof.* By cut-vertex reduction $\text{mr}(F, G) = \sum_{i=1}^h \text{mr}(F, G_i - v) + \min\{2, \sum_{i=1}^k r_v(F, G_i)\}$.
 283 Since $r_v(F, G_1) = 0$, $\text{mr}(F, G) = \text{mr}(F, G_1 - v) + \sum_{i=2}^k \text{mr}(F, G_i - v) + \min\{2, \sum_{i=2}^k r_v(F, G_i)\} =$
 284 $\text{mr}(F, G_1) + \text{mr}(F, G - W_1)$. \square

285 PROPOSITION 2.14. Let $G = (V, E)$ be a graph containing a cycle C on $k \geq 3$ vertices that
286 contains exactly one high degree vertex, v . Then $\text{mr}(F, G) = \text{mr}(F, C) + \text{mr}(F, G - V(C - v))$, or
287 equivalently, $M(F, G) = M(F, G - V(C - v)) + 1$. Furthermore, $Z(G) \leq Z(G - V(C - v)) + 1$. If
288 $M(F, G - V(C - v)) = Z(G - V(C - v))$, then $M(F, G) = Z(G)$.

289 *Proof.* From Proposition 2.13, $\text{mr}(F, G) = \text{mr}(F, C) + \text{mr}(F, G - V(C - v))$, so

$$290 \quad |G| - M(F, G) = (k - 2) + |G| - (k - 1) - M(F, G - V(C - v)),$$

291 or $M(F, G) = M(F, G - V(C - v)) + 1$. To establish $Z(G) \leq Z(G - V(C - v)) + 1$, we exhibit
292 a zero forcing set of order $Z(G - V(C - v)) + 1$. Let B be a minimum zero forcing set for
293 $G - V(C - v)$, and let x be a neighbor of v in C . Then $B \cup \{x\}$ is a zero forcing set for G . If
294 $M(F, G - V(C - v)) = Z(G - V(C - v))$, then $Z(G - V(C - v)) + 1 = M(F, G - V(C - v)) + 1 =$
295 $M(F, G) \leq Z(G) \leq Z(G - V(C - v)) + 1$ so we have equality throughout. \square

296 REMARK 2.15. Every cycle on an even number of vertices is special. Specifically, for a cycle
297 C on $2k$ vertices, the adjacency matrix is optimal if k is even, and if k is odd, an optimal matrix
298 is $A = [a_{ij}] \in \mathcal{S}(F, C)$ where $a_{i, i+1} = 1, i = 1, \dots, 2k - 1$ and $a_{1, 2k} = -1$ (this is valid over every
299 field F).

300 THEOREM 2.16. If G is a graph in \mathcal{K} that does not have a cut-edge, then G is special and
301 $M(F, G) = Z(G)$.

302 *Proof.* We prove the following two statements by induction on the number of cycles for a
303 connected graph $G \in \mathcal{K}$ that does not have a cut-edge.

304 (A) G is a cycle or G contains a cycle with exactly one high degree vertex or G has a subdivided
305 chordal path.

306 (B) G is special and $M(F, G) = Z(G)$.

307 Both (A) and (B) are clear for all cycles in \mathcal{K} , and thus for all connected graphs $G \in \mathcal{K}$ such that
308 G has no cut edge and at most one cycle. Assume both (A) and (B) are true for all connected
309 graphs G having no cut-edge and at most $k \geq 1$ cycles. Let G' be a connected graph in \mathcal{K} that
310 does not have a cut-edge and has $k + 1$ cycles.

311 **Case 1.** G' has a cut-vertex: If G' has a cycle with exactly one high degree vertex, then (A)
312 is true and (B) follows from Proposition 2.14 and the induction hypothesis. If G' does not have a
313 cycle with exactly one high degree vertex, then consider the blocks G_1, \dots, G_b of G' . Since G' has
314 a cut-vertex and no cut-edge, $b > 1$ and each block contains a cycle. Thus G_1 has fewer than $k + 1$
315 cycles. Since G' does not contain a cycle with exactly one high degree vertex, G_1 is not a cycle
316 and does not contain a cycle with at most one high degree vertex. By the induction hypothesis, G_1
317 contains a subdivided chordal path. Since G_1 is a block of G' , G' contains a subdivided chordal
318 path. Thus (A) is true, and (B) follows from Theorem 2.10 and the induction hypothesis.

319 **Case 2.** G' does not have a cut-vertex: Since G' has more than one cycle and G' does not
320 have a cut-vertex, G' has two cycles that intersect in one path on at least two vertices or that
321 intersect in more than one path. Then by Lemma 2.11, G' has two cycles that intersect in one
322 path on at least two vertices. Since $G' \in \mathcal{K}$, by Lemma 2.12, G' has a subdivided chordal path, so
323 (A) is true. Statement (B) then follows from Theorem 2.10 and the induction hypothesis.

324 Since the parameters M and Z sum over connected components, the result for every $G \in \mathcal{K}$
 325 that does not have a cut-edge follows from the result for connected graphs. \square

326 Since \mathcal{K} includes all complete subdivision graphs of simple graphs and multigraphs, we have
 327 the following corollary.

328 **COROLLARY 2.17.** *If G is a simple graph or multigraph that does not have a cut-edge, then*
 329 $M(F, \overline{\overline{G}}) = Z(\overline{\overline{G}})$.

330 **3. Zero forcing number and maximum nullity of edge subdivision graphs.** Recall
 331 that in [3], the authors ask the following question: Suppose G is any graph in which each vertex has
 332 degree at least 3 and H is a graph that has one less edge subdivision than $\overline{\overline{G}}$. Is it always the case
 333 that $M(H) < M(\overline{\overline{G}})$? The graphs G and H given in Example 3.1 below provide a negative answer
 334 to this question. We use the following well known observation: If $G = \cup_{i=1}^h G_i$, $G_i = (V_i, E_i)$, and
 335 (F is infinite or $E_i \cap E_j = \emptyset$ for $i \neq j$), then $\text{mr}(F, G) \leq \sum_{i=1}^h \text{mr}(F, G_i)$.

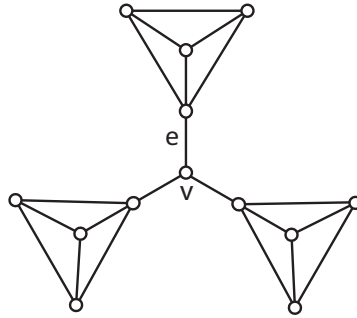


Fig. 3.1: A graph G that provides negative answer to Question 1.7.

336 **EXAMPLE 3.1.** Let G be the graph in Figure 3.1, which is the connected union of three copies
 337 of K_4 (the complete graph on four vertices) and the star graph $K_{1,3}$, with these graphs having no
 338 common edges and the copies of K_4 disjoint; the edge e is one of the edges of the $K_{1,3}$. Let H be
 339 the graph that has one less edge subdivision than $\overline{\overline{G}}$ where the edge e in G is the only unsubdivided
 340 edge. The graphs $\overline{\overline{G}}$ and H are shown in Figure 3.2.

341 Since K_4 has a Hamiltonian path, by Theorem 1.9, $\text{mr}(F, \overline{\overline{K_4}}) = 6$. The subgraph $K_{1,3}$ is a
 342 tree. Hence, by Theorem 1.4, $M(F, \overline{\overline{K_{1,3}}}) = P(\overline{\overline{K_{1,3}}}) = 2$, so $\text{mr}(F, \overline{\overline{K_{1,3}}}) = 5$. Let L be the graph
 343 obtained from $K_{1,3}$ by subdividing all but one edge; again by Theorem 1.4, $M(L) = P(L) = 2$ and
 344 so $\text{mr}(F, L) = 4$. Since $\overline{\overline{G}}$ is a union of three copies of $\overline{\overline{K_4}}$ and one copy of $\overline{\overline{K_{1,3}}}$,

345
$$\text{mr}(F, \overline{\overline{G}}) \leq 3 \text{mr}(F, \overline{\overline{K_4}}) + \text{mr}(F, \overline{\overline{K_{1,3}}}) = 23 \text{ and } M(F, \overline{\overline{G}}) \geq 34 - 23 = 11.$$

346 Similarly, H is a union of three copies of $\overline{\overline{K_4}}$ and one copy of L so

347
$$\text{mr}(F, H) \leq 3 \text{mr}(F, \overline{\overline{K_4}}) + \text{mr}(F, L) = 22 \text{ and } M(F, H) \geq 33 - 22 = 11.$$

348 Furthermore, zero forcing sets of order 11 for both $\overline{\overline{G}}$ and H are exhibited in Figure 3.2. Therefore,
 349 $M(F, H) = Z(H) = M(F, \overline{\overline{G}}) = Z(\overline{\overline{G}}) = 11$.

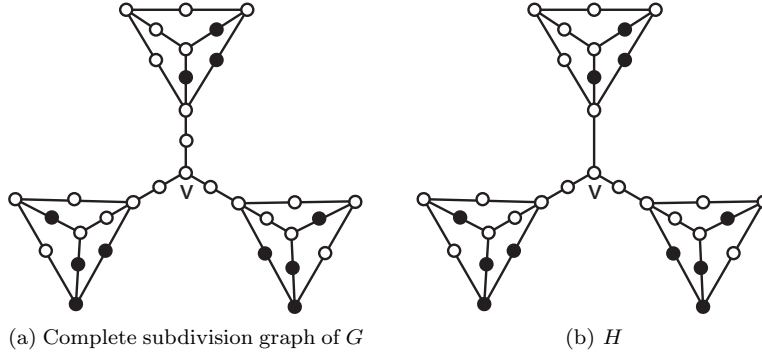


Fig. 3.2: The complete subdivision graph of G and the graph H .

350 Given that we conjecture $M(F, \widetilde{G}) = Z(\widetilde{G})$ for every field F and graph G , one might be
 351 tempted to think that subdividing an edge cannot increase the difference $Z(G) - M(F, G)$. The
 352 next example shows that this is not the case. In fact, $M(F, G) = Z(G)$ does not necessarily imply
 353 $M(F, G_e) = Z(G_e)$.

354 EXAMPLE 3.2. The pentasun H_5 is a five cycle with a degree one neighbor attached to each
 355 cycle vertex, shown in Figure 3.3(a). The graph G in Figure 3.3(b) is obtained from H_5 by adding
 356 two degree one neighbors of u , where u is a vertex of degree one in H_5 . Note the labeled edge $e = uv$;
 357 the result G_e of subdividing edge e is shown in Figure 3.3(c). We show that $M(F, G) = Z(G)$ but
 358 $M(F, G_e) < Z(G_e)$.

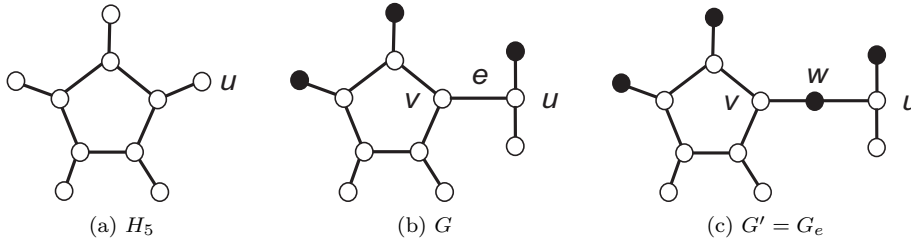


Fig. 3.3: The graphs for Example 3.2

359 It is well known that $M(F, H_5) = 2$, $M(F, H_5 - u) = 2$, $Z(H_5) = 3$, and $Z(H_5 - u) = 2$. Let
 360 $G' := G_e$. The maximum nullity of G and G' can be obtained by performing cut-vertex reduction
 361 using vertex v . Let W_1 (respectively, W'_1) be the vertices in the component of $G - v$ (respectively,
 362 G') containing u and let W_2 (respectively, W'_2) be the vertices of the other component. For $i = 1, 2$,
 363 let $G_i = G[W_i \cup \{v\}]$ and $G'_i = G'[W'_i \cup \{v\}]$. So, $\text{mr}(F, G_1) = 2$, $\text{mr}(F, G[W_1]) = 2$, $\text{mr}(F, G_2) = 7$,
 364 $\text{mr}(F, G[W_2]) = 6$, $\text{mr}(F, G'_1) = 3$, $\text{mr}(F, G'[W'_1]) = 2$, $\text{mr}(F, G'_2) = 7$, and $\text{mr}(F, G'[W'_2]) = 6$.
 365 Thus,

$$366 \quad \text{mr}(F, G) = \sum_{i=1}^2 \text{mr}(F, G[W_i]) + \min\{2, \sum_{i=1}^2 r_v(F, G_i)\} = 9 \text{ so } M(F, G) = 12 - 9 = 3$$

367 and

$$368 \text{mr}(F, G') = \sum_{i=1}^2 \text{mr}(F, G'[W'_i]) + \min\{2, \sum_{i=1}^2 r_v(F, G'_i)\} = 10 \text{ so } M(F, G_e) = M(F, G') = 13 - 10 = 3.$$

369 Zero forcing sets of size 3 for G and 4 for G_e are exhibited in Figures 3.3(b) and 3.3(c), and it
 370 is not difficult to see that no smaller sets can force. Thus $M(F, G) = Z(G) = 3$ and $M(F, G_e) =$
 371 $3 < Z(G_e) = 4$. Zero forcing number and maximum nullity can also be computed by the minimum
 372 rank software [4].

373 It is easy two see that there is no relationship between the change in maximum nullity and
 374 the change in zero forcing number of G and G_e . In Example 3.2 edge subdivision increased zero
 375 forcing number but not maximum nullity. Subdividing any cycle edge of the pentasun H_5 increases
 376 maximum nullity but not zero forcing number (this follows from Proposition 2.1).

377 4. Path cover number of edge subdivision graphs.

378 In this section we investigate the effects of edge subdivisions on the path cover number.

379 PROPOSITION 4.1. *Let G be a graph and e an edge of G . Then*

$$380 P(G) \leq P(G_e) \leq P(G) + 1.$$

381 *If there exists a minimum path cover \mathcal{P} of G such that e is on a path in \mathcal{P} , then $P(G_e) = P(G)$.*

382 *Proof.* Let $e = uv$ and let w be the new vertex in G_e that is adjacent to u and v . We first
 383 prove the upper bounds. Let $\mathcal{P} = \{P_1, \dots, P_k\}$ be a minimum path cover of G . If e is in a path
 384 $Q = P_i$ for some $i = 1 \dots k$, then $(\mathcal{P} \setminus \{Q\}) \cup \{Q_e\}$ is a path cover of G_e , and so $P(G_e) \leq P(G)$. If
 385 e is not in any P_i , then $\mathcal{P} \cup \{w\}$ is a path cover of G_e . In either case, $P(G_e) \leq P(G) + 1$.

386 To prove the lower bound on $P(G_e)$, let $\mathcal{P} = \{P_1, \dots, P_k\}$ be a minimum path cover of G_e .
 387 Then $w \in P_i$ for some i . If $\{w\} = P_i$, then $\mathcal{P} \setminus \{P_i\}$ is a path cover of G . If the edges uw and
 388 vw are in P_i , define P'_i to be the path obtained from P_i by removing uw and vw , and then adding
 389 the edge uv . Then $(\mathcal{P} \setminus \{P_i\}) \cup \{P'_i\}$ is a path cover of G . If w is an endpoint of $P_i \neq \{w\}$, define
 390 P'_i to be the path P_i with w removed. Then $(\mathcal{P} \setminus \{P_i\}) \cup \{P'_i\}$ is a path cover of G . In all cases,
 391 $P(G) \leq P(G_e)$. \square

392 PROPOSITION 4.2. *Let G be a graph and let e be an edge of G . If e is incident to a vertex of*
 393 *degree at most 2, then $P(G_e) = P(G)$.*

394 *Proof.* By Proposition 4.1, $P(G) \leq P(G_e)$. Now it remains to show that $P(G_e) \leq P(G)$. Let
 395 $e = uv$ and let w be the new vertex that is adjacent to u and v in G_e . Without loss of generality,
 396 let $\deg u \leq 2$. Let $\mathcal{P} = \{P_1, \dots, P_k\}$ be a minimum path cover of G . If e is on some path P_i in \mathcal{P} ,
 397 then by Proposition 4.1, $P(G) = P(G_e)$. If e is not in any P_i , then u is the endpoint of some path
 398 in \mathcal{P} . Without loss of generality, say u is in P_1 , then let P'_1 be the path obtained by adding w to
 399 P_1 . Then $(\mathcal{P} \setminus \{P_1\}) \cup \{P'_1\}$ is a path cover of G_e . In either case, $P(G_e) \leq P(G)$. \square

400 It is conjectured that for all graphs G , $M(F, \vec{G}) = Z(\vec{G})$. The following is an example of a
 401 graph G with $P(\vec{G}) < Z(\vec{G})$.

402 EXAMPLE 4.3. Let G be the graph pictured in Figure 4.1, called a double triangle. Since G
 403 contains a Hamiltonian path, by Theorem 1.9, $Z(\overline{G}) = M(F, \overline{G}) = 3$. However, $P(\overline{G}) = 2$ because
 \overline{G} is not a path and a path cover of order 2 is exhibited in Figure 4.1.

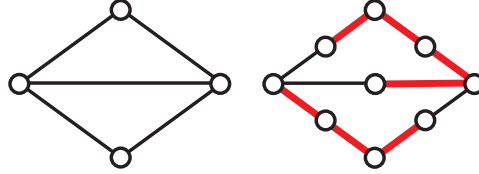


Fig. 4.1: A double triangle and its complete subdivision graph.

404

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