

Minimum Rank and Maximum Eigenvalue Multiplicity of Symmetric Tree Sign Patterns and Trees

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Introduction

Algorithm for Computing Maximum Multiplicity/Minimum Rank

Algorithm

Examples

Rational Matrices Realizing Maximum Multiplicity

Algorithms

Example

Conclusion

- A **graph** G allows loops but does not allow multiple edges.
- A **simple graph** does not allow loops or multiple edges.
- A **tree** T is a connected acyclic graph (ignoring loops).
- A **sign pattern** Z is a matrix with entries in $\{+, -, 0\}$.
- All sign patterns Z and matrices A discussed are **symmetric**, i.e., $z_{ij} = z_{ji}$.
- $\mathcal{G}(Z)$ (or $\mathcal{G}(A)$) is the simple graph with vertices $1, \dots, n$ such that ij is an edge of $\mathcal{G}(Z)$ (or $\mathcal{G}(A)$) if and only if $i \neq j$ and $z_{ij} \neq 0$ ($a_{ij} \neq 0$).
- Z is a **symmetric tree sign pattern** if Z is symmetric and $\mathcal{G}(Z)$ is a tree.
- $\mathcal{S}(G) = \{A : A \text{ is a symmetric matrix and } \mathcal{G}(A) = G\}$.
- $\mathcal{S}(Z) = \{A : A \text{ is a symmetric matrix and for all } i, j, \text{sgn}(a_{ij}) = z_{ij}\}$.
- $S = \mathcal{S}(Z)$ or $\mathcal{S}(T)$.

- The minimum rank of G is defined by

$$\text{mr}(G) = \min_{A \in \mathcal{S}(G)} \{\text{rank}(A)\}.$$

- For A a symmetric matrix, $\text{mult}_\lambda(A)$ denotes the multiplicity of eigenvalue λ of A .
- The maximum multiplicity of an eigenvalue λ in $\mathcal{S}(G)$ is defined by

$$M_\lambda(G) = \max_{A \in \mathcal{S}(G)} \{\text{mult}_\lambda(A)\}.$$

- All nonzero eigenvalues have the same maximum multiplicity.
- $\text{mr}(G) + M_0(G) =$ the number of vertices of G .

- The minimum rank of Z is defined by

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- The maximum multiplicity of an eigenvalue λ in $\mathcal{S}(G)$ is defined by

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- All positive (respectively, negative) eigenvalues have the same maximum multiplicity.
- $\text{mr}(Z) + M_0(Z) = \text{the size of } Z$.

A family S of symmetric matrices **allows** eigenvalue λ if there is a matrix $A \in S$ having eigenvalue λ .

Lemma

- $\mathcal{S}(G)$ allows a nonzero eigenvalue if and only if G has an edge (a loop is an edge).
- $\mathcal{S}(Z)$ allows a positive (respectively, negative) eigenvalue if and only if Z has a nonzero off-diagonal entry or has a positive (negative) diagonal entry.
- Let T be a tree. $\mathcal{S}(T)$ allows 0 eigenvalue if and only if T has no permutation digraphs or at least two permutation digraphs.
- Let Z be a symmetric tree sign pattern. $\mathcal{S}(Z)$ allows eigenvalue 0 if and only if $\det Z$ is identically 0 or has both positive and negative terms.

Algorithm (Computation of maximum multiplicity of λ)

Let T be a simple tree ($\mathcal{G}(Z)$ or a tree with loops suppressed).

Set $Q = \emptyset$ (set of vertices to be deleted).

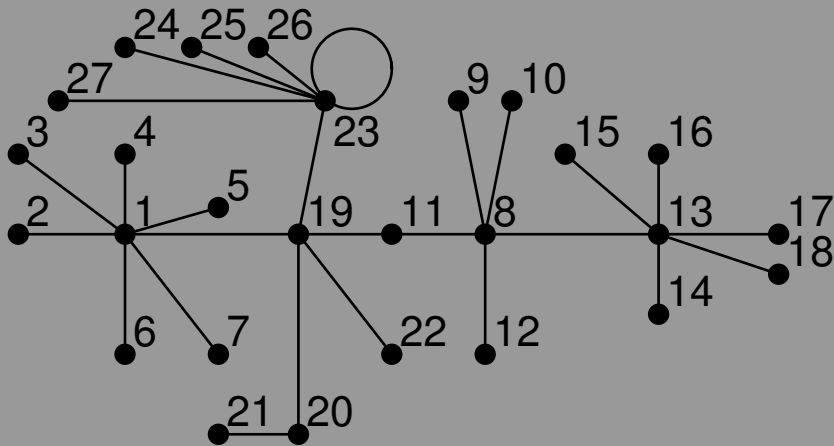
Set $H = \{\text{vertices of } T \text{ with degree } \geq 3\}$ (candidates for deletion).

While $H \neq \emptyset$:

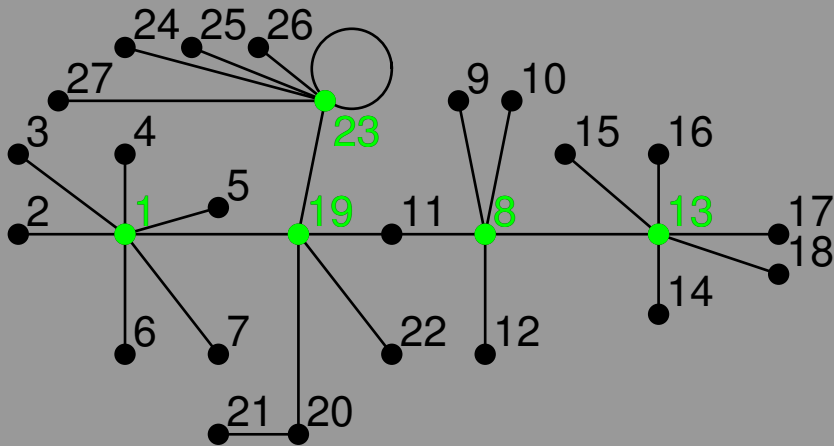
1. Set $\tilde{T} =$ the unique component of $T - Q$ that contains an H -vertex.
2. Set $W = \{w \in H: \text{at most 1 component of } \tilde{T} - w \text{ is not } H\text{-free}\}$.
3. For each $w \in W$, if there are at least two H -free components of $\tilde{T} - w$ that allow eigenvalue λ , then $Q = Q \cup \{w\}$.
4. $H = H - W$.
5. For each $v \in H$, if $\deg_{T-Q} v \leq 2$, remove v from H .

$M_\lambda(Z) = (\text{number of components of } T - Q \text{ that allow } \lambda) - |Q|$.

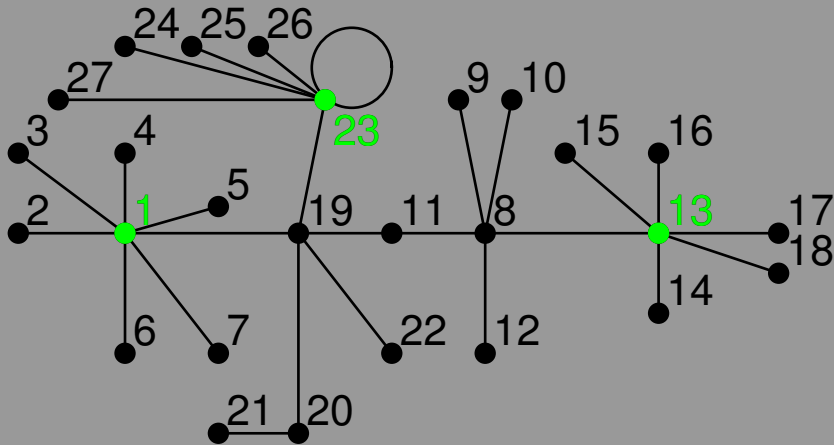
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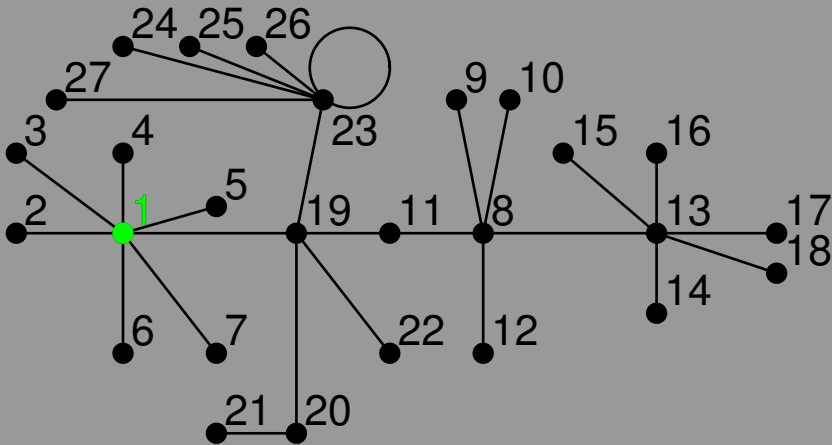


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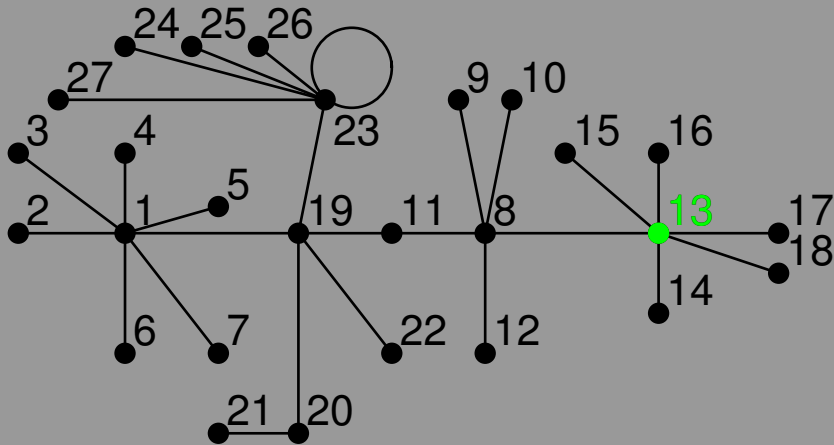
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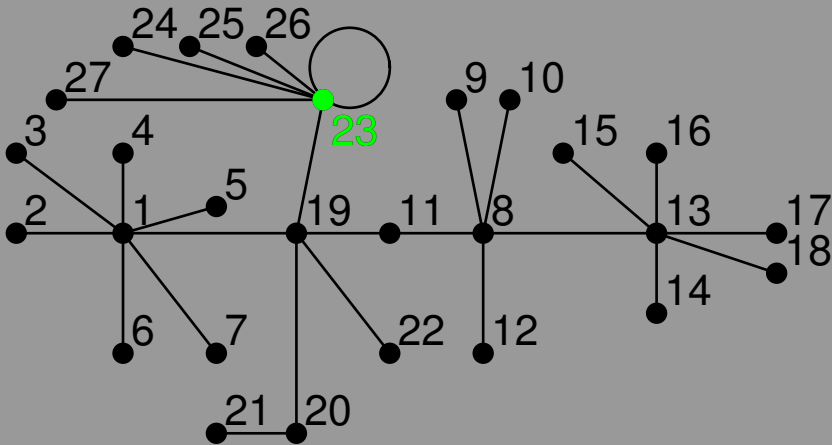
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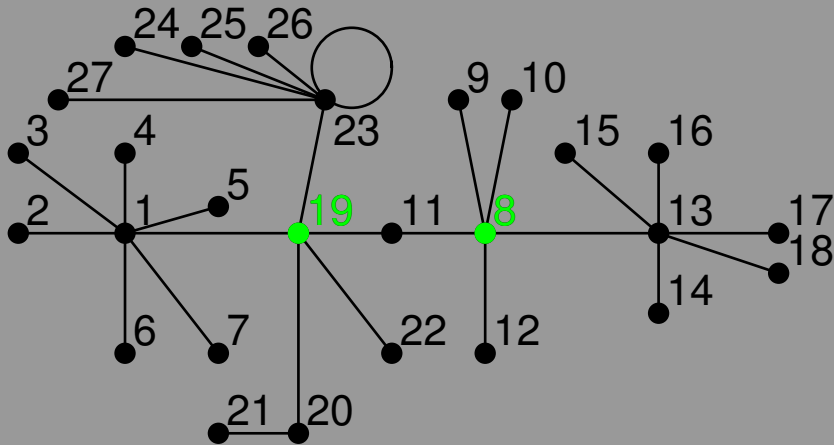
Finish 1st iteration

- 4 $H = H - W$, i.e., $H = \{8, 19\}$.
- 5 For each $v \in H$, if $\deg_{T-Q} v \leq 2$, remove v from H . (no change)

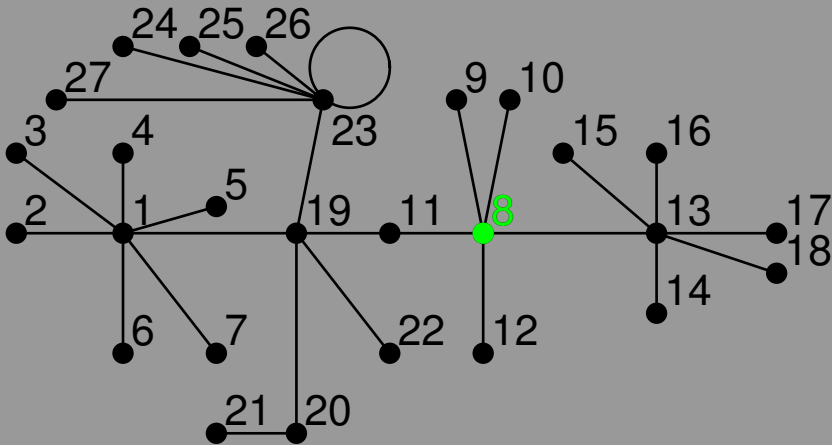
Start 2nd iteration

1. Set $\tilde{T} =$ the unique component of $T - Q$ that contains an H -vertex. (no change)
2. Set $W = \{w \in H: \text{at most 1 component of } \tilde{T} - w \text{ is not } H\text{-free}\}$, i.e., $W = \{8, 19\}$.
3. For each $w \in W$, if there are at least two H -free components of $\tilde{T} - w$ that allow eigenvalue λ , then $Q = Q \cup \{w\}$.

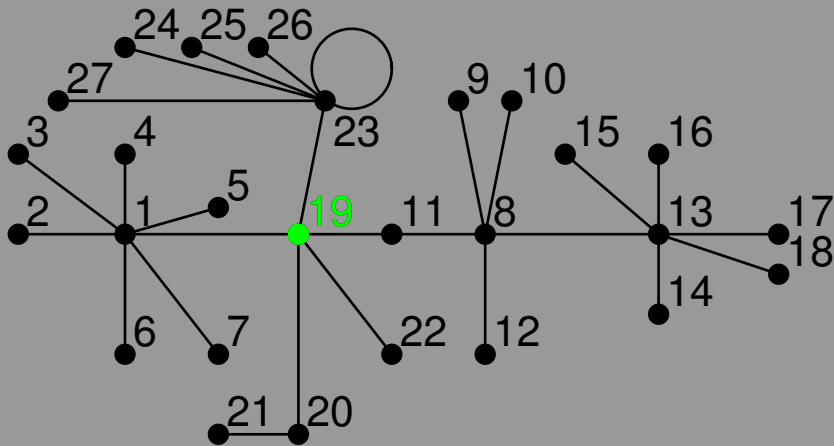
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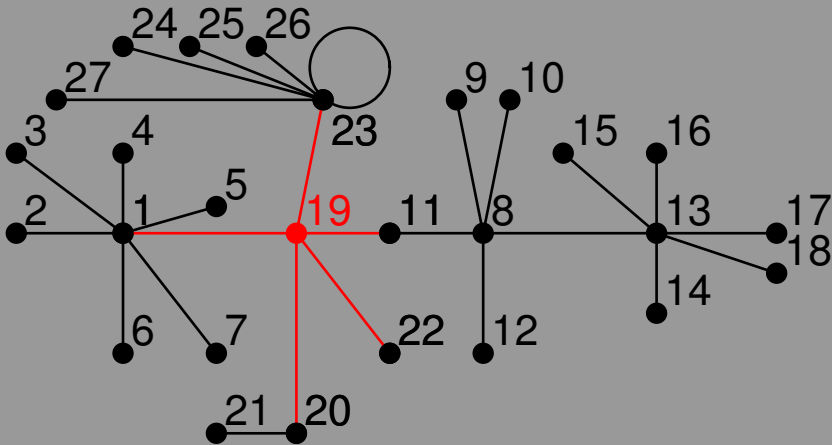
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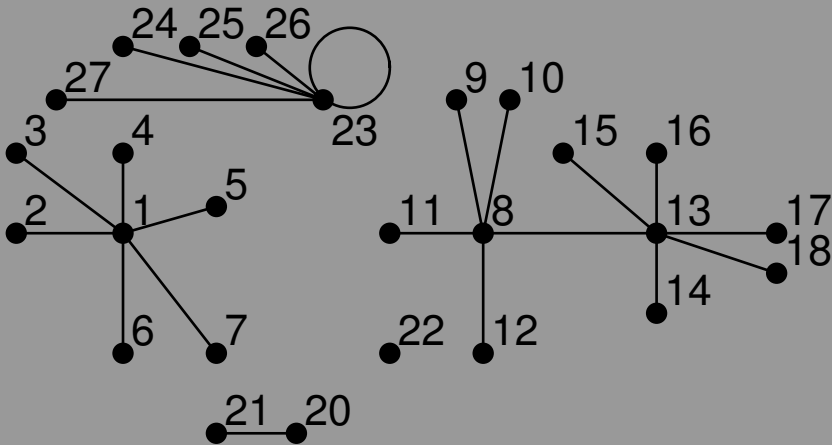
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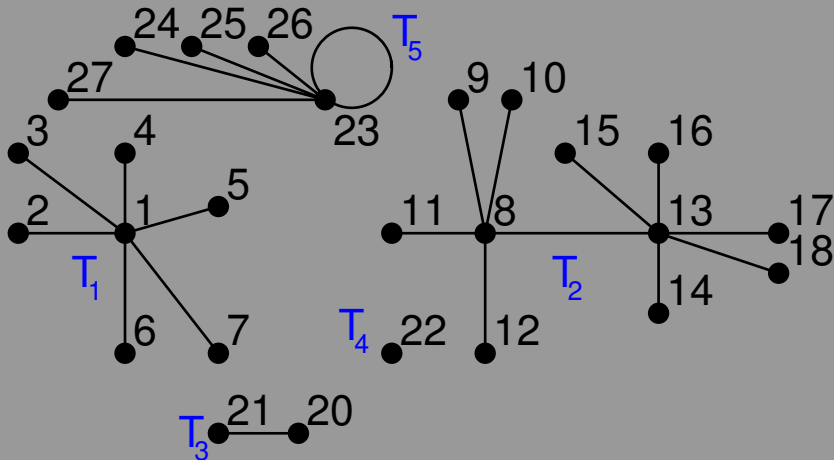
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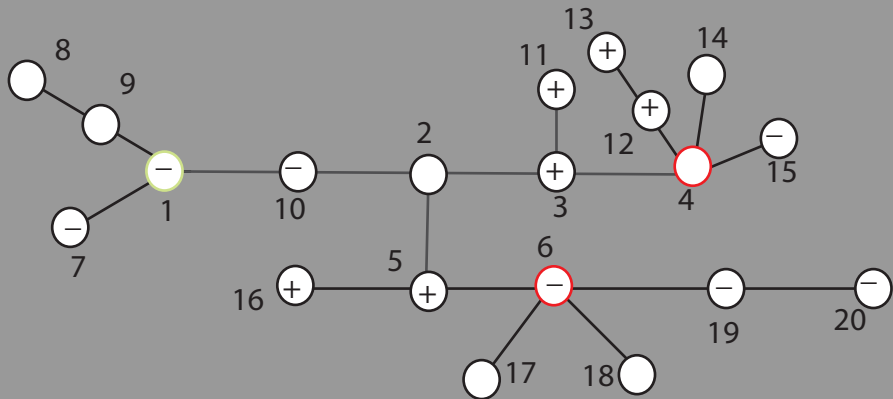


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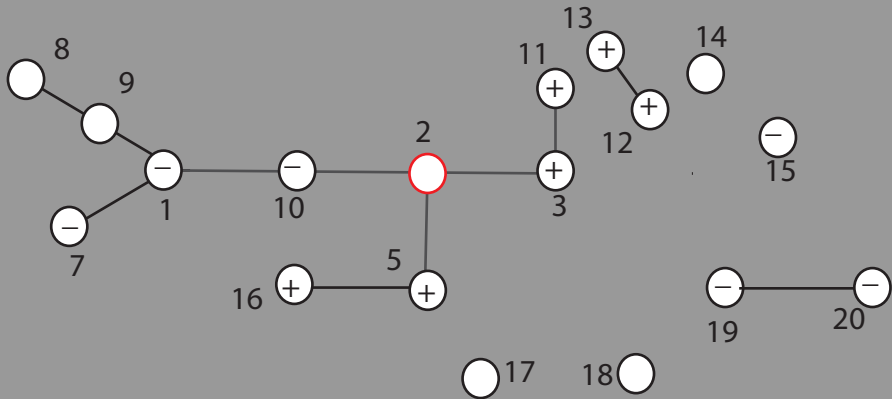


- Example 1: Compute the maximum multiplicity of eigenvalue 1 for the tree:
- $Q = \{19\}$, i.e., 1 vertex was deleted.
- There are 4 components that allow eigenvalue 1, T_1, T_2, T_3, T_5 .
- $M_1(T) = 4 - 1 = 3$.

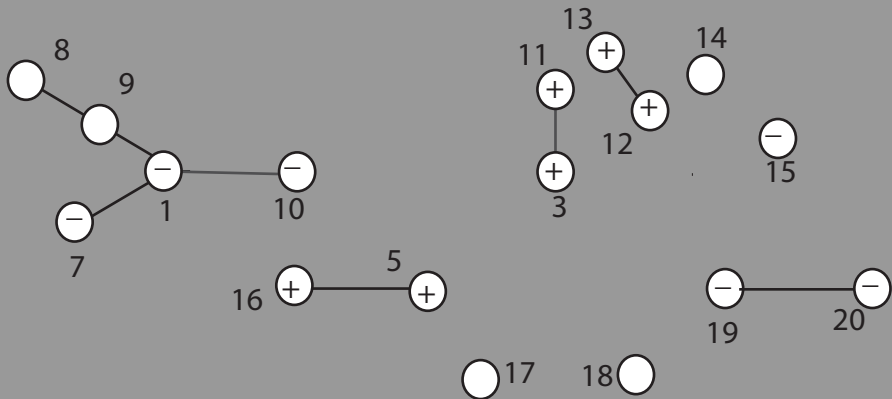
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- $Q = \{2, 4, 6\}$, i.e., 3 vertices were deleted.
- There are 8 components that allow eigenvalue 0.
- $M_0(Z) = 8 - 3 = 5$.
- $\text{mr}(Z) = 20 - 5 = 15$.

The assertions about components allowing eigenvalue 0 rely on results about permutation digraphs.

Algorithm (Construction of a rational matrix realizing maximum multiplicity)

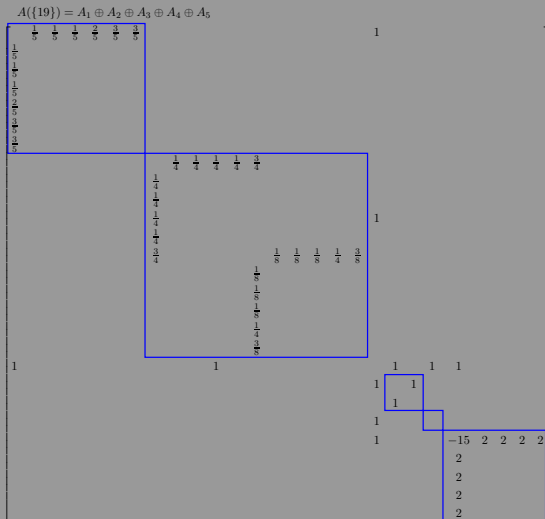
To find a symmetric rational matrix $A \in S = \mathcal{S}(T)$ or $S = \mathcal{S}(Z)$ having maximum multiplicity for rational eigenvalue λ :

- 1. If $\lambda < 0$ and using a sign pattern, replace Z by $-Z$.*
- 2. Apply the algorithm to obtain a set of components.*
- 3. For each component, find a matrix having eigenvalue*
 - 0 if $\lambda = 0$.*
 - 1 if λ is nonzero.*
- 4. Embed these matrices as principal submatrices in a large matrix $A \in S$*
- 5. If $\lambda \neq 0$, multiply A by λ .*

We have algorithms to construct a symmetric rational matrix A satisfying:

- $A \in \mathcal{S}(T)$ and 0 is an eigenvalue of A .
- $A \in \mathcal{S}(Z)$ and 0 is an eigenvalue of A .
- $A \in \mathcal{S}(T)$ and 1 is an eigenvalue of A :
 - If T has a loop.
 - If T does not have a loop.
- $A \in \mathcal{S}(Z)$ and 1 is an eigenvalue of A :
 - If Z has a positive diagonal entry.
 - If Z does not have a positive diagonal entry but has a negative diagonal entry.
 - If all the diagonal entries of Z are 0.

- Example 1: A rational matrix having eigenvalue 1 with multiplicity 3 for the given tree:



Theorem

For any symmetric tree sign pattern Z and any rational number λ , there is a rational matrix $A \in \mathcal{S}(Z)$ such that $\text{mult}_\lambda(A) = M_\lambda(Z)$, and such a matrix can be constructed from the algorithms. For $\lambda = 0$, $\text{rank } A = \text{mr}(Z)$.

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For any tree T and any rational number λ , there is a rational matrix $A \in \mathcal{S}(T)$ such that $\text{mult}_\lambda(A) = M_\lambda(T)$, and such a matrix can be constructed from the algorithms. For $\lambda = 0$, $\text{rank } A = \text{mr}(T)$.