

# Average minimum rank of symmetric matrices described by a graph

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Parameters of random graphs

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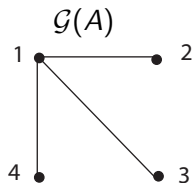
$S_n(\mathbb{R})$  is the set of  $n \times n$  real symmetric matrices.

The graph  $\mathcal{G}(A) = (V, E)$  of  $A \in S_n(\mathbb{R})$  is

- ▶  $V = \{1, \dots, n\}$ ,
- ▶  $E = \{ij : a_{ij} \neq 0 \text{ and } i \neq j\}$ .
- ▶ Diagonal of  $A$  is ignored.

Example:

$$A = \begin{bmatrix} 2 & -1 & 3 & 5 \\ -1 & 0 & 0 & 0 \\ 3 & 0 & -3 & 0 \\ 5 & 0 & 0 & 0 \end{bmatrix}$$



The family of symmetric matrices described by a graph is

$$\mathcal{S}(G) = \{A \in S_n(\mathbb{R}) : \mathcal{G}(A) = G\}.$$

The minimum rank of graph  $G$ :

$$\text{mr}(G) = \min_{A \in \mathcal{S}(G)} \text{rank } A.$$

**Problem** Determine the minimum rank of a graph  $G$ .

## Examples:

Path:  $\text{mr}(P_n) = n - 1$ .Complete graph:  $\text{mr}(K_n) = 1$ .

$$A = \begin{bmatrix} ? & * & 0 & \dots & 0 & 0 \\ * & ? & * & \dots & 0 & 0 \\ 0 & * & ? & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & ? & * \\ 0 & 0 & 0 & \dots & * & ? \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 1 & \dots & 1 \\ 1 & 1 & \dots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \dots & 1 \end{bmatrix}$$

\* is nonzero, ? is indefinite

## Basic properties of minimum rank

- ▶ It is easy to obtain a matrix  $A \in \mathcal{S}(G)$  with  $\text{rank } A = |G| - 1$  (translate).
- ▶ If  $G$  is the disjoint union of graphs  $G_i$  then

$$\text{mr}(G) = \sum \text{mr}(G_i)$$

- ▶ Only connected graphs are studied.
- ▶ If  $G$  is connected,
  - ▶  $\text{mr}(G) = 0$  iff  $G$  is single vertex.
  - ▶  $\text{mr}(G) = 1$  iff  $G = K_n$ ,  $n \geq 2$ .
- ▶  $\text{mr}(G) = |G| - 1$  if and only if  $G$  is a path. [Fiedler 69]

## Results & techniques

Through the work of many people:

Minimum rank is known for

- ▶ Trees
- ▶ Unicyclic graphs
- ▶  $|G| \leq 7$
- ▶  $\text{mr}(G) \leq 2$
- ▶  $\text{mr}(G) \geq |G| - 2$

Techniques for computing minimum rank

- ▶ cut-vertex reduction
- ▶ join reduction
- ▶ Colin de Verdière parameters
- ▶ zero forcing

# Expected value and average minimum rank

Joint work with Tracy Hall, Ryan Martin, Bryan Shader.

- ▶ **average minimum rank** of graphs of order  $n$ :  
sum over all labeled graphs of order  $n$  of minimum ranks of graphs, divided by number of graphs of order  $n$ ,

$$\text{avemr}(n) = \frac{\sum_{|G|=n} \text{mr}(G)}{2^{\binom{n}{2}}}.$$

- ▶  $G(n, p)$  is the Erdős-Rényi random graph on  $n$  vertices with edge probability  $p$ .
- ▶  $\mathbf{E}[\text{mr}(G(n, p))]$  is the expected value of minimum rank
- ▶  $\text{avemr}(n) = \mathbf{E}[\text{mr}(G(n, 1/2))]$



## Theorem

*For  $n$  sufficiently large,*

$$0.146907n < \text{avemr}(n) < 0.5n + \sqrt{7n \ln n}.$$

## Theorem

*With probability approaching 1 as  $n \rightarrow \infty$ ,*

$$|\text{mr}(G(n, 1/2)) - \text{avemr}(n)| < \sqrt{n \ln \ln n}.$$

Expected value of minimum rank is tightly concentrated about the mean

Theorem (Alon, Spencer 2000)

Let  $p \in (0, 1)$ .

Let  $f$  be a graph invariant such that for any graphs  $G$  and  $H$ , if  $x \in V(G) = V(H)$  and  $G - x = H - x$ , then

$$|f(G) - f(H)| \leq 1.$$

Let  $\mu = \mathbf{E}[f(G(n, p))]$ .

Then, for any  $\beta > 0$ ,

$$\Pr[|f(G(n, p)) - \mu| > \beta\sqrt{n-1}] < 2e^{-\beta^2/2}.$$

Proof uses Azuma's inequality for martingales and the vertex exposure martingale.

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## Corollary

*Let  $p \in (0, 1)$  be fixed.*

$$\Pr[|\text{mr}(G(n, p)) - \mathbf{E}[\text{mr}(G(n, p))]| > 2\sqrt{n \ln \ln n}] \\ < 2 \left( \frac{1}{\ln n} \right)^{1/8}.$$

The following two results are well-known, and can be derived from the Chernoff-Hoeffding bound.

$e(G)$  denotes the number of edges of  $G$ .

### Theorem

Let  $p$  be fixed and let  $G$  be distributed according to  $G(n, p)$ .  
Then,

$$e(G) \leq p \binom{n}{2} + n\sqrt{2 \ln n},$$

with probability at least  $1 - n^{-2}$ .

$$e(G) \geq p \binom{n}{2} - n\sqrt{2 \ln n}$$

with probability at least  $1 - n^{-2}$ .

$\delta(G)$  denotes the minimum degree and  $\Delta(G)$  denotes the maximum degree of  $G$ .

## Theorem

Let  $p$  be fixed and let  $G$  be distributed according to  $G(n, p)$ .

Then,

$$pn - \sqrt{6n \ln n} \leq \delta(G) \leq \Delta(G) \leq pn + \sqrt{6n \ln n}$$

with probability at least  $1 - 2n^{-2}$ .

## Lower bound for expected minimum rank

- ▶ **zero-pattern**  $\zeta(\mathbf{x})$  of the real vector  $\mathbf{x} = (x_1, \dots, x_\ell)$ :  $(0, *)$ -vector obtained from  $\mathbf{x}$  by replacing its nonzero entries by  $*$ .
- ▶ **support** of  $\mathbf{z} = (z_1, \dots, z_\ell)$ :  
 $S(\mathbf{z}) = \{i : z_i \neq 0\}$ .

The following result is a variant of a theorem of [Rónyai, Babai, Ganapathy 01]

### Theorem

*If  $\mathbf{f}(\mathbf{x}) = (f_1(\mathbf{x}), f_2(\mathbf{x}), \dots, f_m(\mathbf{x}))$  is an  $m$ -tuple of polynomials in  $k$  variables over a field  $F$  with  $m \geq k$ , each  $f_i$  of degree at most  $d$ , then the number of zero-patterns  $\mathbf{z} = \zeta(\mathbf{f}(\mathbf{x}))$  with  $|S(\mathbf{z})| \leq s$  is at most*

$$\binom{k + sd}{k}.$$

- ▶ Every real symmetric  $n \times n$  matrix of rank at most  $r$  can be expressed in the form

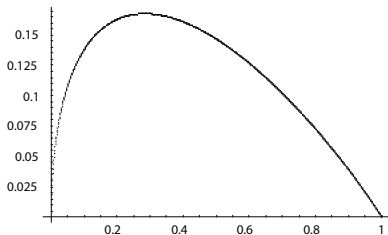
$$X^T D_i X$$

for some  $i$  such that  $0 \leq i \leq r$ , where

- ▶  $D_i = \text{diag}(1, \dots, 1, -1, \dots, -1)$  is an  $r \times r$  diagonal matrix with  $i$  diagonal entries equal to 1 and  $r - i$  equal to  $-1$  and
  - ▶  $X$  is an  $r \times n$  real matrix.
- ▶ Let each entry of  $X$  be a variable.
  - ▶ The total number of variables is  $rn$ .
  - ▶ Each entry of the matrix  $X^T D_i X$  is a polynomial of degree at most 2.

For a fixed value of  $p$  ( $0 < p < 1$ ), let  $c(p)$  be the solution to

$$\frac{(c+p)^{2c+2p}}{(c)^{2c}(p)^{2p}} p^p (1-p)^{(1-p)} = 1.$$



The graph of  $c(p)$

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## Theorem

Let  $G$  be distributed according to  $G(n, p)$  for a fixed  $p$ ,  $0 < p < 1$ . For any  $c < c(p)$ ,

$$\mathbf{E}[\text{mr}(G)] > cn$$

for  $n$  sufficiently large.

## Corollary

For  $n$  sufficiently large, the average minimum rank over all labeled graphs of order  $n$  satisfies

$$\text{avemr}(n) > 0.146907n.$$

## An upper bound for expected minimum rank

- ▶ Recall delta conjecture:  $\text{mr}(G) \leq |G| - \delta(G)$ .
- ▶ Could use this and known bound on  $\delta(G)$  to bound  $\text{mr}(G)$  above.
- ▶ Replace delta conjecture by analogous result for **vertex connectivity**:

$\kappa(G)$  = the smallest number  $k$  such that there is a set of vertices  $S$ , with  $|S| = k$ , for which  $G - S$  is disconnected ( $G \leq K_n$ ).

By convention,  $\kappa(K_n) = n - 1$ .

Terminology taken from [Lovász, Saks, Schrijver 89] – nonstandard for minimum rank problems.

- ▶ **orthogonal representation of  $G$  of dimension  $d$ :**  
 $\varphi : V(G) \rightarrow \mathbb{R}^d$  such that if  $v \not\sim w$ , then  $\varphi(v)$  and  $\varphi(w)$  are orthogonal.
- ▶ Zero representation is orthogonal.
- ▶ **faithful orthogonal representation of  $G$  of dimension  $d$ :**  
 $\varphi : V(G) \rightarrow \mathbb{R}^d$  such that  $v \not\sim w$  if and only if  $\varphi(v)$  and  $\varphi(w)$  are orthogonal.
- ▶ Every faithful orthogonal representation of dimension  $d$  gives rise to a positive semidefinite matrix of rank  $d$  and vice versa.

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## Theorem (Lovász, Saks, Schrijver 89)

*A graph  $G$  has a faithful orthogonal representation of dimension  $|G| - \kappa(G)$ .*

## Corollary

$$\text{mr}(G) \leq |G| - \kappa(G).$$

## Theorem

Let  $G$  be distributed according to  $G(n, p)$ .

- ▶ For  $n$  sufficiently large, the expected value of minimum rank satisfies  $\mathbf{E}[\text{mr}(G)] \leq (1 - p)n + \sqrt{7n \ln n}$ .
- ▶ For  $n$  sufficiently large, the average minimum rank over all labeled graphs of order  $n$  satisfies

$$\text{avemr}(n) \leq 0.5n + \sqrt{7n \ln n}.$$

The proof uses

- ▶  $\text{mr}(G) \leq |G| - \kappa(G)$ ,
- ▶ with probability at least  $1 - n^{-2}$

$$e(G(n, p)) \leq p \binom{n}{2} + n\sqrt{2 \ln n},$$

- ▶  $\delta \leq \frac{2e(G)}{n}$ ,
- ▶ and the relationship (on average) between the connectivity  $\kappa(G)$  and the minimum degree  $\delta(G)$ .

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Bollobás, and Bollobás and Thomason proved that for  $G \sim G(n, p)$ , regardless of  $p$ , then  $\Pr[\kappa(G) < \delta(G)] \rightarrow 0$  as  $n \rightarrow \infty$ .

The following simpler result is obtained by similar means.

### Lemma

*Let  $p \in (0, 1)$  be fixed and  $G$  be distributed according to  $G(n, p)$ . If  $n$  is sufficiently large, then*

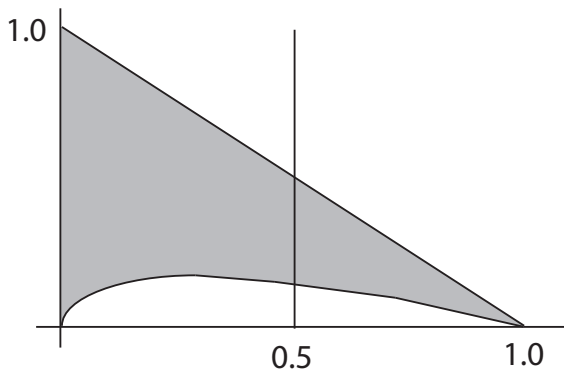
$$\Pr[\kappa(G) < \delta(G)] \leq 3n^{-2}.$$

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The bounds for  $\frac{\mathbf{E}[\text{mr}(G(n,p))]}{n}$  as a function of  $p$



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Thank you!