

Theorem 2.13. Let \mathbf{A}, \mathbf{B} be Σ -algebras and $h \in \text{Hom}(\mathbf{A}, \mathbf{B})$.

- (i) For every $K \in \text{Sub}(\mathbf{A})$, $h(K) \in \text{Sub}(\mathbf{B})$.
- (ii) For every $L \in \text{Sub}(\mathbf{B})$, $h^{-1}(L) := \{a \in A : h(a) \in L\} \in \text{Sub}(\mathbf{A})$.
- (iii) For every $X \subseteq A$, $h(\text{Sg}^{\mathbf{A}}(X)) \in \text{Sg}^{\mathbf{B}}(h(X))$.

Proof. (i). Let $\sigma \in \Sigma_n$ and $b_1, \dots, b_n \in h(K)$. Choose $a_1, \dots, a_n \in K$ such that $h(a_1) = b_1, \dots, h(a_n) = b_n$. Then $\sigma^{\mathbf{B}}(b_1, \dots, b_n) = \sigma^{\mathbf{B}}(h(a_1), \dots, h(a_n)) = h(\sigma^{\mathbf{A}}(a_1, \dots, a_n)) \in h(K)$.

(ii). Let $a_1, \dots, a_n \in h^{-1}(L)$, i.e., $h(a_1), \dots, h(a_n) \in L$. Then $h(\sigma^{\mathbf{A}}(a_1, \dots, a_n)) = \sigma^{\mathbf{B}}(h(a_1), \dots, h(a_n)) \in L$. So $\sigma^{\mathbf{A}}(a_1, \dots, a_n) \in h^{-1}(L)$.

(iii). $h(X) \subseteq h(\text{Sg}^{\mathbf{A}}(X)) \in \text{Sub}(\mathbf{B})$ by part (i). So $\text{Sg}^{\mathbf{B}}(h(X)) \subseteq h(\text{Sg}^{\mathbf{A}}(X))$. For the reverse inclusion, $X \subseteq h^{-1}(h(X)) \subseteq h^{-1}(\text{Sg}^{\mathbf{B}}(h(X))) \in \text{Sub}(\mathbf{A})$ by part (ii). So $\text{Sg}^{\mathbf{A}}(X) \subseteq h^{-1}(\text{Sg}^{\mathbf{B}}(h(X)))$. \square

$h(\mathbf{A})$ will denote the unique subalgebra of \mathbf{B} with universe $h(A)$ and if $\mathbf{B}' \subseteq \mathbf{B}$, then $h^{-1}(\mathbf{B}')$ is the unique subalgebra of \mathbf{A} with universe $h^{-1}(B')$.

Theorem 2.14. Let $\mathbf{A} = \langle A, f \rangle$ be a finite, cyclic mono-unary algebra with period p and tail length l (see Figure 1). Let $h: \mathbf{A} \twoheadrightarrow \mathbf{A}'$ be an epimorphism. Then \mathbf{A}' is finite, cyclic mono-unary algebra. Let p' be its period and l' its tail length. Then p' divides p and $l' \leq l$.

Proof. By Theorem 2.13(iii), \mathbf{A}' is a cyclic mono-unary algebra, and it is obviously finite. Let $A = \text{Sg}^{\mathbf{A}}(\{a\})$. Then $A' = \text{Sg}^{\mathbf{A}'}(\{h(a)\})$ by Theorem 2.13(iii). By definition p is the smallest $m \in \omega \setminus \{0\}$ such that there is an $n \in \omega$ such that $f^{n+m}(a) = f^n(a)$, and l is the smallest $n \in \omega$ such that $f^{n+p}(a) = f^n(a)$. p' and l' are defined similarly. For every $n \geq l$ and every $q \in \omega$, we have

$$(20) \quad f^{n+qp}(a) = f^{n-l}(f^{l+qp}(a)) = f^{n-l}(\underbrace{f^p(f^p(\dots(f^p(f^l(a))))}_{q}) = f^{n-l}(f^l(a)) = f^n(a).$$

We claim that, for all $n, m \in \omega$ with $m > 0$,

$$\text{if } f^{n+m}(a) = f^n(a) \text{ then } p \text{ divides } m.$$

For every $n' \geq n$, $f^{n'+m}(a) = f^{n'-n}(f^{n+m}(a)) = f^{n'-n}(f^n(a)) = f^{n'}(a)$. So without loss of generalization we assume $n \geq l$. By the division algorithm, $m = qp + r$ with $0 \leq r < p$. Then by (20), $f^{n+r}(a) = f^{n+r+qp}(a) = f^{n+m}(a) = f^n(a)$. By the minimality of p , $r = 0$; so $p \mid m$.

$f^{l+p}(h(a)) = h(f^{l+p}(a)) = h(f^l(a)) = f^l(h(a))$. So by (20) (with \mathbf{A}' , $h(a)$, and p' in place of \mathbf{A} , a , and p , respectively), we get that p' divides p . Furthermore, choose q such that $l + qp \geq l'$. Then $f^{l+qp'}(h(a)) = f^{p'}(f^l(h(a))) = f^{p'}(f^{l+qp}(h(a))) = f^{l+qp+p'}(h(a)) = f^{l+qp}(h(a)) = f^l(h(a))$. So $l' \leq l$ by the minimality of l' . \square

Define the binary relation \preceq on $\text{Alg}(\Sigma)$ by $\mathbf{A} \preceq \mathbf{B}$ (equivalently $\mathbf{B} \succcurlyeq \mathbf{A}$) if \mathbf{A} is a homomorphic image of \mathbf{B} , i.e., there is an epimorphism $h: \mathbf{B} \twoheadrightarrow \mathbf{A}$. \preceq is clearly reflexive and it is also transitive, for if $h: \mathbf{B} \twoheadrightarrow \mathbf{A}$ and $g: \mathbf{C} \twoheadrightarrow \mathbf{B}$, then $h \circ g: \mathbf{C} \twoheadrightarrow \mathbf{A}$. However, \preceq fails to be antisymmetric in a strong way: in general,

$$\mathbf{A} \leq \mathbf{B} \text{ and } \mathbf{B} \leq \mathbf{C} \text{ does not imply } \mathbf{A} \cong \mathbf{B}.$$

For example, let $\mathbf{A} = \langle [0, 3], \leq \rangle$ and $\mathbf{B} = \langle [0, 1] \cup [2, 3], \leq \rangle$. Define

$$h(x) = \begin{cases} x & \text{if } 0 \leq x \leq 1, \\ 2 & \text{if } 1 < x < 2, \\ x & \text{if } 2 \leq x \leq 3 \end{cases} \quad \text{and} \quad g(x) = \begin{cases} 3x & \text{if } 0 \leq x \leq 1, \\ 3 & \text{if } 1 < x \leq 3. \end{cases}$$

We leave it as an exercise to prove that h is an epimorphism from the lattice \mathbf{A} to \mathbf{B} and that g is an epimorphism in the opposite direction. However, $\mathbf{A} \not\cong \mathbf{B}$. To see this note that an isomorphism preserves compact elements, but \mathbf{A} had only one compact element, 0, while \mathbf{B} has two, 0 and 2.

If \mathbf{A} or \mathbf{B} is finite, then $\mathbf{A} \preceq \mathbf{B}$ and $\mathbf{B} \preceq \mathbf{A}$ implies $\mathbf{A} \cong \mathbf{B}$. Because, $\mathbf{A} \preceq \mathbf{B}$ and $\mathbf{B} \preceq \mathbf{A}$ imply $|A| \leq |B|$ and $|B| \leq |A|$, i.e., $|A| = |B|$. So any surjective homomorphism from \mathbf{A} onto \mathbf{B} must be also injective by the pigeonhole principle. Thus $\mathbf{A} \cong \mathbf{B}$.

\cong is an equivalence relation on $\text{Alg}(\Sigma)$. ($\Delta_A: \mathbf{A} \cong \mathbf{A}$; if $h: \mathbf{A} \cong \mathbf{B}$ and $g: \mathbf{B} \cong \mathbf{C}$, then $g \circ h: \mathbf{A} \cong \mathbf{C}$; if $h: \mathbf{A} \cong \mathbf{B}$ then $h^{-1}: \mathbf{B} \cong \mathbf{A}$.) The equivalence class of $[\mathbf{A}]_{\cong}$ of \mathbf{A} under \cong , which we normally write simply as $[\mathbf{A}]$, is called the *isomorphism type* of \mathbf{A} . ($[\mathbf{A}]$ is not a proper set, it's too big, but this problem can be disregarded for our purposes.) The class of all isomorphism types of Σ -algebras is denoted by $[\text{Alg}(\Sigma)]$.

The relations of subalgebra and homomorphic image on $\text{Alg}(\Sigma)$ induce corresponding relations on $[\text{Alg}(\Sigma)]$.

- $[\mathbf{A}] \subseteq [\mathbf{B}]$ if $\mathbf{A} \cong ; \subseteq \mathbf{B}$, i.e., if $\exists \mathbf{C} (\mathbf{A} \cong \mathbf{C} \subseteq \mathbf{B})$.
- $[\mathbf{A}] \preceq [\mathbf{B}]$ if $\mathbf{A} \preceq \mathbf{B}$. (Note that because $\cong \subseteq \preceq$, $(\cong ; \preceq) \subseteq (\preceq ; \preceq) = \preceq$.)

\subseteq and \preceq are well defined in isomorphism types, i.e., if $\mathbf{A} \cong \mathbf{A}'$ and $\mathbf{B} \cong \mathbf{B}'$, then $\mathbf{A} \cong ; \subseteq \mathbf{B}$ iff $\mathbf{A}' \cong ; \subseteq \mathbf{B}'$ and $\mathbf{A} \preceq \mathbf{B}$ iff $\mathbf{A}' \preceq \mathbf{B}'$.

To see that these equivalences holds we observe that $\mathbf{A} \cong ; \subseteq \mathbf{B}$ implies $\mathbf{A}' \cong ; \cong ; \subseteq ; \cong \mathbf{B}'$, and $\mathbf{A} \preceq \mathbf{B}$ implies $\mathbf{A} \cong ; \preceq ; \cong \mathbf{B}$. The second implication holds because clearly $\cong ; \preceq = \preceq ; \cong = \preceq$. The first implication is an immediate consequence of the equality $\subseteq ; \cong = \cong ; \subseteq$, which is in turn a corollary of Thm. 2.15(i) below.

\subseteq is a partial ordering of isomorphism types, and \preceq is what is called a *quasi-ordering* or *pre-ordering*, i.e., it is reflexive and transitive but not symmetric. However, \preceq is a partial ordering on *finite isomorphism types*, that is isomorphism types of finite algebras. Clearly, if $[\mathbf{A}] \preceq [\mathbf{B}]$ and $[\mathbf{B}] \preceq [\mathbf{C}]$ and \mathbf{A} (equivalently \mathbf{B}) is finite, then $[\mathbf{A}] = [\mathbf{B}]$.

Let us consider the various relative products of \subseteq and \preceq and their converses:

$$\subseteq ; \preceq, \preceq ; \subseteq, \subseteq ; \succ, \succ ; \subseteq.$$

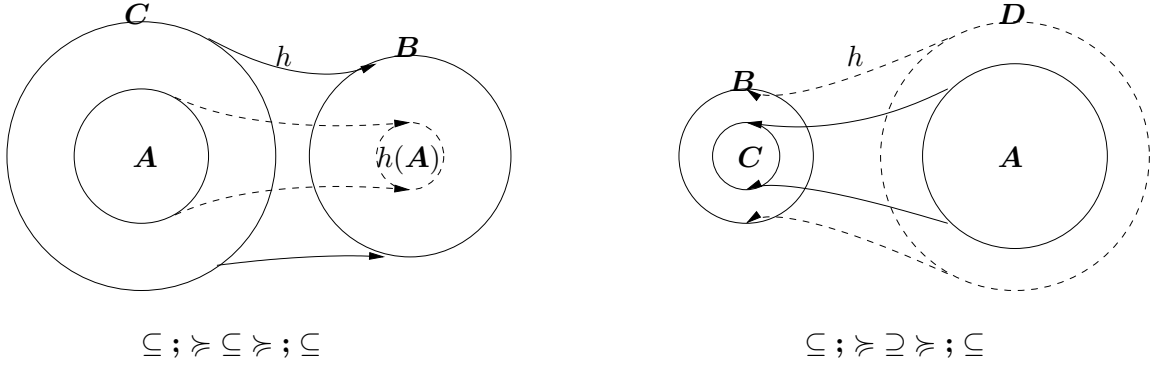
This gives half of the eight possible combinations, but each of the remaining four is a converse of one of these. For example, $(\supseteq ; \preceq) = (\overset{\sim}{\subseteq} ; \overset{\sim}{\preceq}) = (\subseteq ; \preceq)^\sim$.

Theorem 2.15. *The following inclusions as relations be Σ -isomorphism types.*

- (i) $\subseteq ; \succ = \succ ; \subseteq$.
- (ii) $\subseteq ; \preceq \subseteq \preceq ; \subseteq$.

Proof. (i) \subseteq . Assume $[\mathbf{A}] \subseteq ; \succ [\mathbf{B}]$, i.e., there exists a \mathbf{C} such that $\mathbf{A} \subseteq \mathbf{C} \succ \mathbf{B}$. We need to show $[\mathbf{A}] \succ ; \subseteq [\mathbf{B}]$, i.e., there exists a \mathbf{D} such that $\mathbf{A} \succ \mathbf{D} \subseteq \mathbf{B}$. Let $h: \mathbf{C} \rightarrow \mathbf{B}$. Then $\mathbf{A} \succ h(\mathbf{A}) \subseteq \mathbf{B}$. See the following figure.

The inclusion \supseteq of (i) is left as an exercise. See the following figure.



(ii) \subseteq . Assume $\mathbf{A} \subseteq \mathbf{C} \preccurlyeq \mathbf{B}$. Let $h: \mathbf{B} \rightarrow \mathbf{C}$. Then $\mathbf{A} \preccurlyeq h^{-1}(\mathbf{A}) \subseteq \mathbf{B}$.

We show by example that the inclusion of (ii) is proper. Let $\mathbf{Q} = \langle \mathbb{Q}, +, \cdot, -, 0, 1 \rangle$ be the ring of rational numbers (a field). Recall that $\mathbf{Z} = \langle \mathbb{Z}, +, \cdot, -, 0, 1 \rangle$ is the ring of integers and let $\mathbf{Z}_2 = \langle \mathbb{Z}_2, +, \cdot, -, 0, 1 \rangle$ be the ring of integers (mod 2). We know that $\mathbf{Z}_2 \preccurlyeq \mathbf{Z} \subseteq \mathbf{Q}$, so $\mathbf{Z}_2 \preccurlyeq ; \subseteq \mathbf{Q}$. But it is not the case that $\mathbf{Z}_2 \subseteq ; \preccurlyeq \mathbf{Q}$. In fact, we show that

$$(21) \quad \mathbf{H}(\{\mathbf{Q}\}) = \mathbf{I}(\{\mathbf{Q}\}) \cup \{ \mathbf{A} \in \text{Alg}(\Sigma) : |\mathbf{A}| = 1 \},$$

i.e., the only nontrivial (two or more elements) homomorphic images of \mathbf{Q} are its isomorphic images. Suppose $h: \mathbf{Q} \rightarrow \mathbf{A}$, and suppose h is not an isomorphism, i.e., it is not injective. Let a and b be distinct elements of \mathbf{Q} such that $h(a) = h(b)$. $a - b \neq 0$ but $h(a - b) = h(a + -b) = h(a) + -h(b) = h(a) - h(b) = 0$. Thus $1 = h(1) = h((a - b) \cdot (a - b)^{-1}) = h(a - b) \cdot h((a - b)^{-1}) = 0 \cdot h((a - b)^{-1}) = 0$. So for every $a \in \mathbf{A}$, $a = 1 \cdot a = 0 \cdot a = 0$; i.e., \mathbf{A} is trivial. This proves the claim.

Suppose now by way of contradiction that for some \mathbf{A} , $\mathbf{Z}_2 \subseteq \mathbf{A} \preccurlyeq \mathbf{Q}$. By the claim \mathbf{A} must be either isomorphic to \mathbf{Q} or a trivial one-element algebra. But \mathbf{Z}_2 is not isomorphic to a subalgebra of \mathbf{Q} . \square

For any class \mathbf{K} of Σ -algebras, we define

$$\mathbf{H}(\mathbf{K}) = \{ \mathbf{A} \in \text{Alg}(\Sigma) : \exists \mathbf{B} \in \mathbf{K} (\mathbf{A} \preccurlyeq \mathbf{B}) \},$$

$$\mathbf{I}(\mathbf{K}) = \{ \mathbf{A} \in \text{Alg}(\Sigma) : \exists \mathbf{B} \in \mathbf{K} (\mathbf{A} \cong \mathbf{B}) \},$$

the classes respectively of homomorphic and isomorphic images of algebras of \mathbf{K} . \mathbf{H} and \mathbf{I} are algebraic closure operators on $\text{Alg}(\Sigma)$. For example $\mathbf{H}\mathbf{H}(\mathbf{K}) = \mathbf{H}(\mathbf{K})$ because of the transitivity of \preccurlyeq . \mathbf{H} is algebraic because $\mathbf{H}(\mathbf{K}) = \bigcup \{ \mathbf{H}(\mathbf{A}) : \mathbf{A} \in \mathbf{K} \}$.

Theorem 2.16. For any class \mathbf{K} of Σ -algebras,

(i) $\mathbf{S}\mathbf{H}(\mathbf{K}) \subseteq \mathbf{H}\mathbf{S}(\mathbf{K})$.

(ii) $\mathbf{H}\mathbf{S}$ is an algebraic closure operator on $\text{Alg}(\Sigma)$.

Proof. (i). Suppose $\mathbf{A} \in \mathbf{S}\mathbf{H}(\mathbf{K})$. Then there exists a $\mathbf{B} \in \mathbf{K}$ such that $\mathbf{A} \subseteq ; \preccurlyeq \mathbf{B}$. Then by Thm. 2.15(ii), $\mathbf{A} \preccurlyeq ; \subseteq \mathbf{B}$. Thus $\mathbf{A} \in \mathbf{H}\mathbf{S}(\mathbf{K})$.

(ii) $\mathbf{K} \subseteq \mathbf{S}(\mathbf{K})$ by the extensivity of \mathbf{K} , and hence by the extensivity and monotonicity of \mathbf{H} , $\mathbf{K} \subseteq \mathbf{H}(\mathbf{K}) \subseteq \mathbf{H}\mathbf{S}(\mathbf{K})$. So $\mathbf{H}\mathbf{S}$ is extensive. $\mathbf{H}\mathbf{S}\mathbf{H}\mathbf{S}(\mathbf{K}) \stackrel{(i)}{\subseteq} \mathbf{H}\mathbf{H}\mathbf{S}\mathbf{S}(\mathbf{K}) = \mathbf{H}\mathbf{S}(\mathbf{K})$. Since clearly $\mathbf{H}\mathbf{S}\mathbf{H}\mathbf{S}(\mathbf{K}) \subseteq \mathbf{H}\mathbf{S}(\mathbf{K})$, we get that $\mathbf{H}\mathbf{S}$ is idempotent. Finally, $\mathbf{K} \subseteq \mathbf{L}$ implies $\mathbf{S}(\mathbf{K}) \subseteq \mathbf{S}(\mathbf{L})$ which in turn implies $\mathbf{H}\mathbf{S}(\mathbf{K}) \subseteq \mathbf{H}\mathbf{S}(\mathbf{L})$. So $\mathbf{H}\mathbf{S}$ is monotonic. $\mathbf{A} \in \mathbf{H}\mathbf{S}(\mathbf{K})$ iff

there is a $\mathbf{B} \in \mathbf{K}$ such that $\mathbf{A} \preccurlyeq ; \subseteq \mathbf{B}$. Thus $\mathbf{HS}(\mathbf{K}) \subseteq \bigcup \{ \mathbf{HS}(\mathbf{B}) : \mathbf{B} \in \mathbf{K} \}$. Thus \mathbf{HS} is algebraic. \square

From Thm. 2.15(i) we see that the opposite inclusion of Thm. 2.16(i) does not hold in general.

We also note the following obvious identities, which prove useful later. $\mathbf{I}(\mathbf{K}) = \mathbf{HI}(\mathbf{K}) = \mathbf{H}(\mathbf{K})$ and $\mathbf{IS}(\mathbf{K}) = \mathbf{SI}(\mathbf{K})$.

Let Σ be a multi-sorted signature with sort set S . Let \mathbf{A} and \mathbf{B} be Σ -algebras. A *homomorphism* $h: \mathbf{A} \rightarrow \mathbf{B}$ is a S -sorted map $h = \langle h_s : s \in S \rangle$ such that, for all $s \in S$, $h_s: A_s \rightarrow B_s$, and for all $\sigma \in \Sigma$ of type $s_1, \dots, s_n \rightarrow s$ and for all $\langle a_1, \dots, a_n \rangle \in A_{s_1} \times \dots \times A_{s_n}$, $h_s(\sigma^{\mathbf{A}}(a_1, \dots, a_n)) = \sigma^{\mathbf{B}}(h_{s_1}(a_1), \dots, h_{s_n}(a_n))$.

Example. Let A and B be nonempty sets. Recall that

$$\mathbf{Lists}(A) = \langle \langle A \cup \{e_D\}, A^* \cup \{e_l\} \rangle, \text{head}, \text{tail}, \text{append}, \text{emptylist}, \text{derror}, \text{derror} \rangle.$$

Let $f: A \rightarrow B$ be any map. We define the S -sorted map $h = \langle h_D, h_L \rangle$ where $h_D \upharpoonright A = f$ and $h_D(e_D) = e_D$, and, for all $a_1, \dots, a_n \in D$, $h_L(\langle a_1, \dots, a_n \rangle) = \langle f(a_1), \dots, f(a_n) \rangle$, and $h_L(e_l) = e_l$. Then $h \in \text{Hom}(\mathbf{A}, \mathbf{B})$ and every homomorphism from \mathbf{A} to \mathbf{B} comes from some $f: A \rightarrow B$ in this way (Exercise).

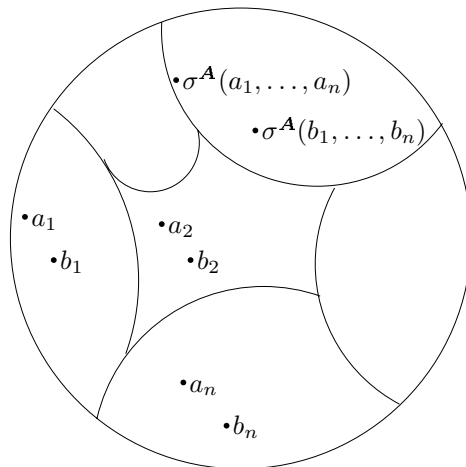
2.5. Congruence relations and quotient algebras.

Definition 2.17. Let \mathbf{A} be a Σ -algebra. An equivalence relation E on A is called a *congruence relation* if, for all $n \in \omega$, all $\sigma \in \Sigma_n$, and all $a_1, \dots, a_n, b_1, \dots, b_n \in A$,

$$(22) \quad a_1 E b_1, \dots, a_n E b_n \text{ imply } \sigma^{\mathbf{A}}(a_1, \dots, a_n) E \sigma^{\mathbf{A}}(b_1, \dots, b_n).$$

The set of all congruences on \mathbf{A} is denoted by $\text{Co}(\mathbf{A})$.

(22) is called the *substitution property*. Intuitively, it asserts that the equivalence class of the result of applying any one of the fundamental operations of \mathbf{A} depends only on the equivalence classes of the arguments. See the following figure.



We use lower case Greek letters, e.g., α, β, γ , etc., to represent congruence letters. The equivalence class $[a]_\alpha$ of a is called the *congruence class* of a and is normally denoted by

a/α . So $a/\alpha \sim b/\alpha$ iff $a/\alpha = b/\alpha$ iff $a \in b/\alpha$. The set of all congruence classes of α , i.e., the partition of \mathbf{A} , is denoted by \mathbf{A}/α .

Definition 2.18. Let \mathbf{A} be a Σ -algebra and let $\alpha \in \text{Co}(\mathbf{A})$. We define an Σ -algebra

$$\mathbf{A}/\alpha = \langle \mathbf{A}/\alpha, \sigma^{\mathbf{A}/\alpha} \rangle_{\sigma \in \Sigma}$$

on the set of congruence classes of α as follows. For every $n \in \omega$, every $\sigma \in \Sigma_n$, and for all $a_1/\alpha, \dots, a_n/\alpha \in \mathbf{A}/\alpha$,

$$\sigma^{\mathbf{A}/\alpha}(a_1/\alpha, \dots, a_n/\alpha) = \sigma^{\mathbf{A}}(a_1, \dots, a_n)/\alpha.$$

\mathbf{A}/α is called a *quotient algebra*, or more precisely, the *quotient of \mathbf{A} by α* .

Note that $\sigma^{\mathbf{A}/\alpha}$ is well defined by the substitution property.

Examples. (1) $\Delta_A, \nabla_A \in \text{Co}(\mathbf{A})$. $\mathbf{A}/\Delta_A \cong \mathbf{A}$ and \mathbf{A}/∇_A is a trivial one-element algebra.

(2) Let $\mathbf{G} = \langle G, \cdot, ^{-1}, e \rangle$ be a group. A subuniverse N of \mathbf{G} is called *normal* if it is closed under *conjugation* by arbitrary elements of G , i.e., if $a \in N$ implies $x \cdot a \cdot x^{-1} \in N$ for every $x \in G$. Define $a \sim b$ if $a \cdot b^{-1} \in N$. Then \sim is a congruence relation. Furthermore, for each congruence α on \mathbf{G} , e/α is a normal subuniverse of \mathbf{G} . The mapping $\alpha \mapsto e/\alpha$ is a bijection between $\text{Co}(\mathbf{G})$ and the normal subuniverses of \mathbf{G} . We leave the verification of all these facts to the reader but will give the details of a similar verification for rings.

(3) Let $\mathbf{R} = \langle R, +, \cdot, -, 0 \rangle$ be a ring, and let I be an *ideal* subuniverse of \mathbf{R} : if $a \in I$ then $x \cdot a, a \cdot x \in I$ for every $x \in R$. Define $a \sim b$ if $a - b (= a + -b)$ is in I . We show \sim is an equivalence relation. $a - a = 0 \in I$. $a - b, b - c \in I$ imply $a - c = (a - b) + (b - c) \in I$. $a - b \in I$ implies $b - a = -(a - b) \in I$. We now verify the substitution property. Suppose $a_1 \sim b_1$ and $a_2 \sim b_2$, i.e., $a_1 - b_1, a_2 - b_2 \in I$. Then $(a_1 + a_2) - (b_1 + b_2) = (a_1 - b_1) + (a_2 - b_2) \in I$. So $(a_1 + a_2) \sim (b_1 + b_2)$. $-a_1 - (-b_1) = -(a_1 - b_1) \in I$. So $-a_1 \sim -b_1$. Finally, $((a_1 - b_1) \cdot a_2) + (b_1 \cdot (a_2 - b_2)) \in I$. But $((a_1 - b_1) \cdot a_2) + (b_1 \cdot (a_2 - b_2)) = a_1 \cdot a_2 - b_1 \cdot a_2 + b_1 \cdot a_2 - b_1 \cdot b_2 = a_1 \cdot a_2 - b_1 \cdot b_2$. So $(a_1 \cdot a_2) \sim (b_1 \cdot b_2)$. Thus \sim is a congruence on \mathbf{R} .

Now let $\alpha \in \text{Co}(\mathbf{R})$. We check that $0/\alpha$ is an ideal. Let $a, a' \in 0/\alpha$, i.e., $a \sim a' \sim 0$. Then $(a + a') \sim (0 + 0) = 0$, $-a \sim -0 = 0$, $(a \cdot a') \sim (0 \cdot 0) = 0$. So $0/\alpha$ is a subuniverse. Moreover, for every $x \in R$, $(x \cdot a) \sim (x \cdot 0) = 0$ and $(a \cdot x) \sim (0 \cdot x) = 0$. So $0/\alpha$ is an ideal. Furthermore, $(a - b) \in 0/\alpha$ iff $(a - b) \sim 0$ iff $a = ((a - b) + b) \sim (0 + b) = b$. So \sim is the congruence determined as in the first part by the ideal $0/\alpha$. Conversely, if one starts with an ideal I , constructs the congruence \sim as in the first part, then forms its ideal $0/\sim$ one gets back I , because $a \in 0/\sim$ iff $a \sim 0$ iff $a = (a - 0) \in I$. So for any ring \mathbf{R} the mapping $\alpha \rightarrow 0/\alpha$ is a bijection between $\text{Co}(\mathbf{R})$ and the set of ideals of \mathbf{R} .